

IRANIAN JOURNAL OF Mathematical Chemistry Journal homepage: ijmc.kashanu.ac.ir



Original Scientific Paper

# On the Graovac–Ghorbani and Atom–Bond Connectivity Indices of Graphs from Primary Subgraphs

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# ARTICLE INFO

# ABSTRACT

Article History: Received: 8 January 2022 Accepted: 27 February 2022 Published online: 30 March 2022 Academic Editor: Mehdi Eliasi

#### **Keywords:**

Atom-bond connectivity index Graovac-Ghorbani index Cactus graphs Let G = (V, E) be a finite simple graph. The Graovac-Ghorbani index of a graph G is defined as

 $ABC_{GG}(G) = \sum_{uv \in E(G)} \sqrt{\frac{n_u(uv,G) + n_v(uv,G) - 2}{n_u(uv,G)n_v(uv,G)}}$ 

where  $n_u(uv, G)$  is the number of vertices closer to vertex u than vertex v of the edge  $uv \in E(G)$ .  $n_v(uv, G)$  is defined analogously. The atom-bond connectivity index of a graph G is defined as

$$ABC(G) = \sum_{uv \in E(G)} \sqrt{\frac{d_u + d_v - 2}{d_u d_v}}$$

where  $d_u$  is the degree of vertex u in G. Let G be a connected graph constructed from pairwise disjoint connected graphs  $G_1, \ldots, G_k$  by selecting a vertex of  $G_1$ , a vertex of  $G_2$ , and identifying these two vertices. Then continue in this manner inductively. We say that G is obtained by point-attaching from  $G_1, \ldots, G_k$  and that  $G_i$ 's are the primary subgraphs of G. In this paper, we give some upper bounds on Graovac-Ghorbani and atom-bond connectivity indices for these graphs. Additionally, we consider some particular cases of these graphs that are of importance in chemistry and study their Graovac-Ghorbani and atom-bond connectivity indices.

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# **1. INTRODUCTION**

A molecular graph is a simple graph such that its vertices correspond to the atoms and the edges to the bonds of a molecule. Let G = (V, E) be a finite, connected, simple graph. A topological index of G is a real number related to G. It does not depend on the labeling or

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DOI: 10.22052/IJMC.2022.246079.1612

pictorial representation of a graph. The Wiener index W(G) is the first distance based topological index defined as  $W(G) = \sum_{\{u,v\} \subseteq G} d(u,v) = \frac{1}{2} \sum_{u,v \in V(G)} d(u,v)$  with the summation runs over all pairs of vertices of *G* [26]. The topological indices and graph invariants based on distances between vertices of a graph are widely used for characterizing molecular graphs, establishing relationships between structure and properties of molecules, predicting biological activity of chemical compounds, and making their chemical applications. The Wiener index is one of the most used topological indices with high correlation with many physical and chemical indices of molecular compounds [26]. In 2010, Graovac et al. [14] introduced a new bond-additive structural invariant as a quantitative refinement of the distance nonbalancedness and also a measure of peripherality in graphs. They used the name Graovac-Ghorbani index for this invariant which is defined as

$$ABC_{GG}(G) = \sum_{uv \in E(G)} \sqrt{\frac{n_u(uv,G) + n_v(uv,G) - 2}{n_u(uv,G)n_v(uv,G)}}$$

where  $n_u(uv, G)$  is the number of vertices of G closer to u than to v, and similarly,  $n_v(uv, G)$  is the number of vertices closer to v than to u. Equidistant vertices from u and v are not taken into account to compute  $n_u(uv, G)$  and  $n_v(uv, G)$ . They determined some bounds on this index. Graovac et al. in [15] computed that for some nanostar dendrimers. Some other upper and lower bounds on the  $ABC_{GG}$  index and also characterizing the extremal graphs was studied by Das [4]. Ghorbani et al. in [13] calculated the  $ABC_{GG}$  of an infinite family of fullerenes. More results on this index can be found in [5, 10, 20, 22, 23].

Graovac and Ghorbani defined  $ABC_{GG}(G)$  [14] which motivated by the definition of atom-bond connectivity index. Initially, the atom-bond connectivity index of a graph G, ABC(G), was defined [9] as:

$$ABC(G) = \sqrt{2} \sum_{uv \in E(G)} \sqrt{\frac{d_u + d_v - 2}{d_u d_v}},$$

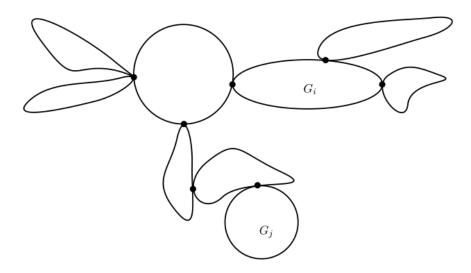
but later on, this index was very slightly redefined [8] by dropping the factor  $\sqrt{2}$ . We refer the reader to [1] for a complete review of the atom-bond connectivity index.

Cactus graphs were first known as Husimi tree, they appeared in the scientific literature more than sixty years ago in papers by Husimi and Riddell concerned with cluster integrals in the theory of condensation in statistical mechanics [16, 18, 21]. We refer the reader to [2, 3, 11, 12, 17, 24, 25] for some aspects of parameters of cactus graphs.

In this paper, we consider the Graovac-Ghorbani and atom-bond connectivity indices of graphs from primary subgraphs. For convenience, the definition of these kind of graphs will be given in the next section. In Section 2, we obtain some upper bounds for Graovac-Ghorbani and atom-bond connectivity indices of graphs from primary subgraphs. In Section 3, we obtain the Graovac-Ghorbani and atom-bond connectivity indices of families of graphs that are of importance in chemistry.

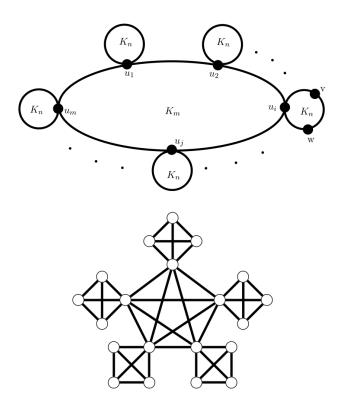
# 2. MAIN RESULTS

Let G be a connected graph constructed from pairwise disjoint connected graphs  $G_1, ..., G_k$  as follows. Select a vertex of  $G_1$ , a vertex of  $G_2$ , and identify these two vertices. Then continue in this manner inductively. Note that the graph G constructed in this way has a tree-like structure, the  $G_i$ 's being its building stones, see Figure 1.



**Figure 1**: A graph with subgraph units  $G_1, \ldots, G_k$ .

Usually say that G is obtained by point-attaching from  $G_1, ..., G_k$  and that  $G_i$ 's are the primary subgraphs of G. A particular case of this construction is the decomposition of a connected graph into blocks (see [7]). We consider some particular cases of these graphs and study their atom-bond connectivity index. As an example of point-attaching graph, consider the graph  $K_m$  and m copies of  $K_n$ . By definition, the graph Q(m,n) is obtained by identifying each vertex of  $K_m$  with a vertex of a unique  $K_n$ . The graph Q(5,4)is shown in Figure 2.



**Figure 2**: The graph Q(m, n) and Q(5,4), respectively.

**Theorem 2.1.** For the graph Q(m, n) (see Figure 2), and n ≥ 2 we have:  
i. 
$$ABC(Q(m, n)) = \frac{m(m-1)}{2(m+n-2)} \sqrt{2(m+n-3)} + m(\frac{n}{2}-1)\sqrt{2(n-2)}$$
  
 $+ m(n-1) \sqrt{\frac{m+2n-5}{n^2+mn-m-3n+2}}.$   
ii.  $ABC_{GG}(Q(m, n)) = \frac{m(m-1)}{2n} \sqrt{2n-2} + m(n-1) \sqrt{\frac{n(m-1)}{n(m-1)+1}}.$ 

Proof.

(i) There are 
$$\frac{m(m-1)}{2}$$
 edges with endpoints of degree  $m + n - 2$ . Also there are  $m(n-1)$  edges with endpoints of degree  $m + n - 2$  and  $n - 1$ , and there are  $m(n-1)(\frac{n}{2}-1)$  edges with endpoints of degree  $n - 1$ . Therefore

$$ABC(Q(m,n)) = \frac{m(m-1)}{2} \sqrt{\frac{(m+n-2)+(m+n-2)-2}{(m+n-2)(m+n-2)}} + m(n-1) \sqrt{\frac{(m+n-2)+(n-1)-2}{(m+n-2)(n-1)}} + m(n-1)(\frac{n}{2}-1) \sqrt{\frac{(n-1)+(n-1)-2}{(n-1)(n-1)}},$$

and we have the result.

(ii) First consider the edge  $u_i u_j$  in  $K_m$ . There are *n* vertices which are closer to  $u_i$  than  $u_j$  (including  $u_i$  itself), also there are *n* vertices closer to  $u_j$  than  $u_i$ , and there are  $\frac{m(m-1)}{2}$  edges like  $u_i u_j$  in Q(m, n). Now consider the edge vw in the *i*-th  $K_n$ . There is one vertex which is closer to *v* than *w* and that is *v* itself, and visa versa. Finally, consider the edge  $u_i v$  in the *i*-th  $K_n$ . There are n(m-1) + 1 vertices which are closer to  $u_i$  than *v* (including  $u_i$ ), also there is one vertex closer to *v* than  $u_i$  which is *v*, and there are m(n-1) edges like  $u_i v$  in Q(m, n).

Therefore we have the result.

## 2.1 UPPER BOUNDS

By the definition of the atom-bond connectivity and Graovac-Ghorbani indices, we have the following easy result:

**Proposition 2.2.** Let G be a disconnected graph with components  $G_1$  and  $G_2$  Then

- i.  $ABC(G) = ABC(G_1) + ABC(G_2)$ .
- ii.  $ABC_{GG}(G) = ABC_{GG}(G_1) + ABC_{GG}(G_2).$

Now we examine the effects on ABC(G) and  $ABC_{GG}(G)$  when G is modified by deleting an edge or vertex of G.

**Theorem 2.3.** Let G = (V, E) be a graph and  $e = uv \in E$  which is not a pendant edge. Also let  $d_u$  be the degree of vertex u in G, and  $n_u$  be the number of vertices of G closer to u than to v. Then,

i. 
$$ABC(G - e) \ge ABC(G) - \max\left\{\frac{\sqrt{2d_u - 2}}{d_v}, \frac{\sqrt{2d_v - 2}}{d_u}\right\}.$$
  
ii.  $ABC_{GG}(G - e) \ge ABC_{GG}(G) - \max\left\{\frac{\sqrt{2n_u - 2}}{n_v}, \frac{\sqrt{2n_v - 2}}{n_u}\right\}$ 

**Proof.** First we remove edge *e* and find ABC(G - e). For every integer  $a, b \ge 2$ , we have

 $\sqrt{\frac{a+(b-1)-2}{a(b-1)}} \ge \sqrt{\frac{a+b-2}{ab}}$ . Now Obviously, by adding edge *e* to G - e and  $\sqrt{\frac{d_u+d_v-2}{d_ud_v}}$  to ABC(G - e), then ABC(G) is less than that or equal to it. So

$$ABC(G) \le ABC(G-e) + \sqrt{\frac{d_u + d_v - 2}{d_u d_v}}$$
$$\le ABC(G-e) + \max\left\{\sqrt{\frac{d_u + d_u - 2}{d_v d_v}}, \sqrt{\frac{d_v + d_v - 2}{d_u d_u}}\right\}$$

$$= ABC(G-e) + \max\left\{\frac{\sqrt{2d_u-2}}{d_v}, \frac{\sqrt{2d_v-2}}{d_u}\right\}$$

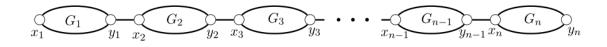
and therefore we have the result. The proof is similar to Part (i).

By the same argument as the proof of Theorem 2.3, and deleting a vertex at the first step, we have:

**Theorem 2.4.** Let G = (V, E) be a graph and  $v \in V$ . Also let  $d_u$  be the degree of vertex u in G. Then,

i. 
$$ABC(G - v) \ge ABC(G) - \sum_{uv \in E} \max\left\{\frac{\sqrt{2d_u - 2}}{d_v}, \frac{\sqrt{2d_v - 2}}{d_u}\right\}.$$
  
ii.  $ABC_{GG}(G - v) \ge ABC_{GG}(G) - \sum_{uv \in E} \max\left\{\frac{\sqrt{2n_u - 2}}{n_v}, \frac{\sqrt{2n_v - 2}}{n_u}\right\}.$ 

Here we study some bounds on the atom-bond connectivity and Graovac-Ghorbani indices for links of graphs and circuits of graphs.



**Figure 3**: Link of *n* graphs  $G_1, G_2, \ldots, G_n$ .

**Theorem 2.5.** Let  $G_1, G_2, ..., G_k$  be a finite sequence of pairwise disjoint connected graphs and let  $x_i, y_i \in V(G_i)$ . Let G be the link of graphs  $\{G_i\}_{i=1}^k$  with respect to the vertices  $\{x_i, y_i\}_{i=1}^k$ , Figure 3, and suppose that  $G_i \neq K_1$ . Then,

i. 
$$ABC(G) \leq \sum_{i=1}^{k} ABC(G_i) + \sum_{i=1}^{k-1} \max\left\{\frac{\sqrt{2d_{x_{i+1}}-2}}{d_{y_i}}, \frac{\sqrt{2d_{y_i}-2}}{d_{x_{i+1}}}\right\}.$$
  
ii.  $ABC_{GG}(G) \leq \sum_{i=1}^{k} ABC_{GG}(G_i) + \sum_{i=1}^{k-1} \max\left\{\frac{\sqrt{2n_{x_{i+1}}-2}}{n_{y_i}}, \frac{\sqrt{2n_{y_i}-2}}{n_{x_{i+1}}}\right\}.$ 

**Proof.** We first remove the edge  $y_1x_2$ , see Figure 3. By Theorem 2.3, we have

$$ABC(G) \le ABC(G - y_1 x_2) + \max\{\frac{\sqrt{2d_{x_2} - 2}}{d_{y_1}}, \frac{\sqrt{2d_{y_1} - 2}}{d_{x_2}}\}.$$

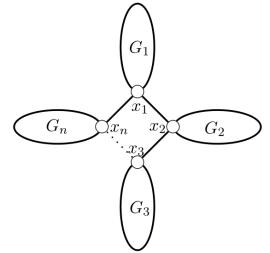
Let G' be the link graph related to graphs  $\{G_i\}_{i=2}^k$  with respect to the vertices  $\{x_i, y_i\}_{i=2}^k$ . Then by Proposition 2.2 we have,

$$ABC(G - y_1x_2) = ABC(G_1) + ABC(G'),$$

and therefore,

$$ABC(G) \le ABC(G_1) + ABC(G') + \max\{\frac{\sqrt{2d_{x_2}-2}}{d_{y_1}}, \frac{\sqrt{2d_{y_1}-2}}{d_{x_2}}\}.$$

By continuing this process, we have the result. The proof of (ii) is similar to Part (i).



**Figure 4**: Circuit of *n* graphs  $G_1, G_2, \ldots, G_n$ .

**Theorem 2.6.** Let  $G_1, G_2, ..., G_k$  be a finite sequence of pairwise disjoint connected graphs and let  $x_i \in V(G_i)$ . Let G be the circuit of graphs  $\{G_i\}_{i=1}^k$  with respect to the vertices  $\{x_i\}_{i=1}^k$  and obtained by identifying the vertex  $x_i$  of the graph  $G_i$  with the *i*-th vertex of the cycle graph  $C_k$  (Figure 4) and suppose that  $G_i \neq K_1$ . Then,

i. 
$$ABC(G) \le \max\left\{\frac{\sqrt{2d_{x_1}-2}}{d_{x_n}}, \frac{\sqrt{2d_{x_n}-2}}{d_{x_1}}\right\} + \sum_{i=1}^k ABC(G_i)$$
  
  $+ \sum_{i=1}^{k-1} \max\left\{\frac{\sqrt{2d_{x_{i+1}}-2}}{d_{x_i}}, \frac{\sqrt{2d_{x_i}-2}}{d_{x_{i+1}}}\right\},$   
ii.  $ABC_{GG}(G) \le \max\left\{\frac{\sqrt{2n_{x_1}-2}}{n_{x_n}}, \frac{\sqrt{2n_{x_n}-2}}{n_{x_1}}\right\} + \sum_{i=1}^k ABC_{GG}(G_i)$   
  $+ \sum_{i=1}^{k-1} \max\left\{\frac{\sqrt{2n_{x_{i+1}}-2}}{n_{x_i}}, \frac{\sqrt{2n_{x_i}-2}}{n_{x_{i+1}}}\right\}.$ 

**Proof.** First we remove the edge  $x_n x_1$ , Figure 4. By Theorem 2.3, we have

$$ABC(G) \le ABC(G - x_n x_1) + \max\{\frac{\sqrt{2d_{x_1} - 2}}{d_{x_n}}, \frac{\sqrt{2d_{x_n} - 2}}{d_{x_1}}\}.$$

Now we remove edge  $x_1x_2$ . Then,

$$ABC(G) \le ABC(G - \{x_n x_1, x_1 x_2\}) + \max\{\frac{\sqrt{2d_{x_1} - 2}}{d_{x_n}}, \frac{\sqrt{2d_{x_n} - 2}}{d_{x_1}}\} + \max\{\frac{\sqrt{2d_{x_2} - 2}}{d_{x_1}}, \frac{\sqrt{2d_{x_1} - 2}}{d_{x_2}}\}.$$

Let G' be the graph related to circuit graph with  $\{G_i\}_{i=2}^k$  with respect to the vertices  $\{x_i\}_{i=2}^k$  and removing the edge  $x_n x_1$ . Then by Proposition 2.2 we have,

$$ABC(G - \{x_n x_1, x_1 x_2\}) = ABC(G_1) + ABC(G'),$$

and therefore,

$$ABC(G) \le ABC(G_1) + ABC(G') + \max\{\frac{\sqrt{2d_{x_1}-2}}{d_{x_n}}, \frac{\sqrt{2d_{x_n}-2}}{d_{x_1}}\} + \max\{\frac{\sqrt{2d_{x_2}-2}}{d_{x_1}}, \frac{\sqrt{2d_{x_1}-2}}{d_{x_2}}\}.$$

By continuing this process, we have the result. The proof of (ii) is similar to Part (i).

## 2.2 Some Other Upper Bounds for the Graovac–Ghorban Iindex

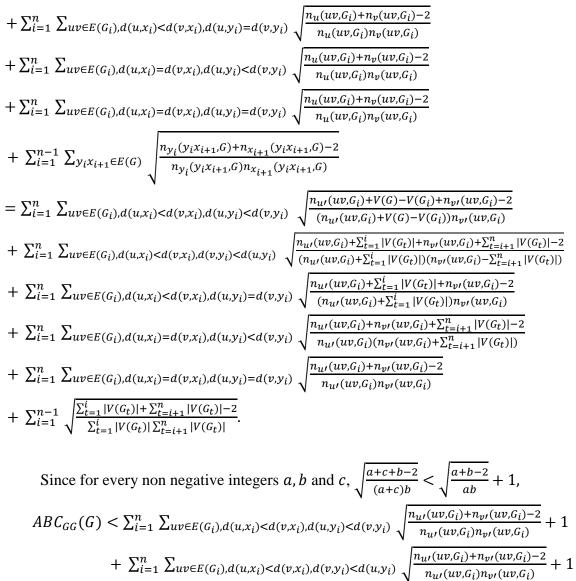
In this subsection, we consider some special graphs from primary subgraphs and present upper bounds for the Graovac-Ghorbani index of them. The following theorem is about the link of graphs.

**Theorem 2.7.** Let  $G_1, G_2, ..., G_k$  be a finite sequence of pairwise disjoint connected graphs and let  $x_i, y_i \in V(G_i)$ . Let G be the link of graphs  $\{G_i\}_{i=1}^k$  with respect to the vertices  $\{x_i, y_i\}_{i=1}^k$ , see Figure 3. Then,

$$ABC_{GG}(G) < (|E(G)| - (n-1)) + \sum_{i=1}^{n} ABC_{GG}(G_i) + \sum_{i=1}^{n-1} \sqrt{\frac{|V(G)| - 2}{\sum_{i=1}^{i} |V(G_i)| \sum_{i=i+1}^{n} |V(G_i)|}}.$$

**Proof.** Consider the graph  $G_i$  (Figure 3) and let  $n_{u'}(uv, G_i)$  be the number of vertices of  $G_i$  closer to u than v in  $G_i$ , Also let  $n_u(uv, G_i)$  be the number of vertices of G closer to u than v in G. By the definition of Graovac-Ghorbani index, we have:

$$\begin{aligned} ABC_{GG}(G) &= \sum_{uv \in E(G)} \sqrt{\frac{n_u(uv,G) + n_v(uv,G) - 2}{n_u(uv,G)n_v(uv,G)}} \\ &= \sum_{i=1}^n \sum_{uv \in E(G_i)} \sqrt{\frac{n_u(uv,G_i) + n_v(uv,G_i) - 2}{n_u(uv,G_i)n_v(uv,G_i)}} \\ &+ \sum_{i=1}^{n-1} \sum_{y_i x_{i+1} \in E(G)} \sqrt{\frac{n_{y_i}(y_i x_{i+1},G) + n_{x_{i+1}}(y_i x_{i+1},G) - 2}{n_{y_i}(y_i x_{i+1},G)n_{x_{i+1}}(y_i x_{i+1},G)}} \\ &= \sum_{i=1}^n \sum_{uv \in E(G_i), d(u,x_i) < d(v,x_i), d(u,y_i) < d(v,y_i)} \sqrt{\frac{n_u(uv,G_i) + n_v(uv,G_i) - 2}{n_u(uv,G_i)n_v(uv,G_i)}}} \\ &+ \sum_{i=1}^n \sum_{uv \in E(G_i), d(u,x_i) < d(v,x_i), d(v,y_i) < d(u,y_i)} \sqrt{\frac{n_u(uv,G_i) + n_v(uv,G_i) - 2}{n_u(uv,G_i)n_v(uv,G_i)}}} \end{aligned}$$



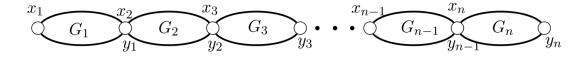
$$\begin{split} C_{GG}(G) &< \sum_{i=1}^{n} \sum_{uv \in E(G_{i}), d(u,x_{i}) < d(v,x_{i}), d(u,y_{i}) < d(v,y_{i})} \sqrt{\frac{n_{u'}(uv,G_{i}) + n_{v'}(uv,G_{i}) - 2}{n_{u'}(uv,G_{i}) n_{v'}(uv,G_{i})}} + 1 \\ &+ \sum_{i=1}^{n} \sum_{uv \in E(G_{i}), d(u,x_{i}) < d(v,x_{i}), d(v,y_{i}) < d(u,y_{i})} \sqrt{\frac{n_{u'}(uv,G_{i}) + n_{v'}(uv,G_{i}) - 2}{n_{u'}(uv,G_{i}) n_{v'}(uv,G_{i})}} + 1 \\ &+ \sum_{i=1}^{n} \sum_{uv \in E(G_{i}), d(u,x_{i}) < d(v,x_{i}), d(u,y_{i}) = d(v,y_{i})} \sqrt{\frac{n_{u'}(uv,G_{i}) + n_{v'}(uv,G_{i}) - 2}{n_{u'}(uv,G_{i}) n_{v'}(uv,G_{i})}} + 1 \\ &+ \sum_{i=1}^{n} \sum_{uv \in E(G_{i}), d(u,x_{i}) = d(v,x_{i}), d(u,y_{i}) < d(v,y_{i})} \sqrt{\frac{n_{u'}(uv,G_{i}) + n_{v'}(uv,G_{i}) - 2}{n_{u'}(uv,G_{i}) n_{v'}(uv,G_{i})}} + 1 \\ &+ \sum_{i=1}^{n} \sum_{uv \in E(G_{i}), d(u,x_{i}) = d(v,x_{i}), d(u,y_{i}) < d(v,y_{i})} \sqrt{\frac{n_{u'}(uv,G_{i}) + n_{v'}(uv,G_{i}) - 2}{n_{u'}(uv,G_{i}) n_{v'}(uv,G_{i})}} + 1 \\ &+ \sum_{i=1}^{n} \sum_{uv \in E(G_{i}), d(u,x_{i}) = d(v,x_{i}), d(u,y_{i}) = d(v,y_{i})} \sqrt{\frac{n_{u'}(uv,G_{i}) + n_{v'}(uv,G_{i}) - 2}{n_{u'}(uv,G_{i}) n_{v'}(uv,G_{i})}} + 1 \\ &+ \sum_{i=1}^{n-1} \sqrt{\frac{\sum_{i=1}^{i} |V(G_{t})| + \sum_{i=i+1}^{n} |V(G_{t})| - 2}{\sum_{i=1}^{i} |V(G_{t})| + \sum_{i=i+1}^{n} |V(G_{t})| - 2}}} \\ &= (|E(G)| - (n - 1)) + \sum_{i=1}^{n} ABC_{GG}(G_{i}) \\ &+ \sum_{i=1}^{n-1} \sqrt{\frac{\sum_{i=1}^{i} |V(G_{t})| + \sum_{i=i+1}^{n} |V(G_{t})| - 2}{\sum_{i=1}^{i} |V(G_{t})| + \sum_{i=i+1}^{n} |V(G_{t})| - 2}}}, \end{split}$$

and therefore we have the result.

By the same argument similar to the proof of the Theorem 2.7, we have the following theorem which is about the chain of graphs:

**Theorem 2.8.** Let  $G_1, G_2, ..., G_n$  be a finite sequence of pairwise disjoint connected graphs and let  $x_i, y_i \in V(G_i)$ . Let  $C(G_1, ..., G_n)$  be the chain of graphs  $\{G_i\}_{i=1}^n$  with respect to the vertices  $\{x_i, y_i\}_{i=1}^k$  which obtained by identifying the vertex  $y_i$  with the vertex  $x_{i+1}$  for i = 1, 2, ..., n - 1 (Figure 5). Then,

 $ABC_{GG}(C(G_1,\ldots,G_n)) < |E(G)| + \sum_{i=1}^n ABC_{GG}(G_i) + \sum_{i=1}^{n-1} \sqrt{\frac{|V(G)| - 2}{\sum_{t=1}^i |V(G_t)| \sum_{t=i+1}^n |V(G_t)|}}.$ 

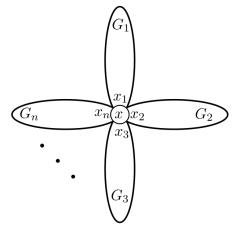


**Figure 5**: Chain of *n* graphs  $G_1, G_2, \ldots, G_n$ .

With similar argument to the proof of the Theorem 2.7, we have the following theorem which is about the bouquet of graphs:

**Theorem 2.9.** Let  $G_1, G_2, ..., G_n$  be a finite sequence of pairwise disjoint connected graphs and let  $x_i \in V(G_i)$ . Let  $B(G_1, ..., G_n)$  be the bouquet of graphs  $\{G_i\}_{i=1}^n$  with respect to the vertices  $\{x_i\}_{i=1}^n$  and obtained by identifying the vertex  $x_i$  of the graph  $G_i$  with x (see Figure 6). Then,

$$ABC_{GG}(B(G_1,\ldots,G_n)) < |E(G)| + \sum_{i=1}^n ABC_{GG}(G_i) + \sum_{i=1}^{n-1} \sqrt{\frac{|V(G)| - 2}{\sum_{i=1}^i |V(G_t)| \sum_{i=i+1}^n |V(G_t)|}}.$$



**Figure 6**: Bouquet of *n* graphs  $G_1, G_2, ..., G_n$  and  $x_1 = x_2 = \cdots = x_n = x$ .

Now we consider the circuit of graphs.

**Theorem 2.10.** Let  $G_1, G_2, ..., G_n$  be a finite sequence of pairwise disjoint connected graphs and let  $x_i \in V(G_i)$ . Let G be the circuit of graphs  $\{G_i\}_{i=1}^n$  with respect to the vertices  $\{x_i\}_{i=1}^n$  and obtained by identifying the vertex  $x_i$  of the graph  $G_i$  with the *i*-th vertex of the cycle graph  $C_n$ , Figure 4. Then,

$$ABC_{GG}(G) < (|E(G)| - n) + \sum_{i=1}^{n} ABC_{GG}(G_i) + \sqrt{\frac{|V(G)| - 2}{|V(G_1)||V(G_n)|}} + \sum_{i=1}^{n-1} \sqrt{\frac{|V(G)| - 2}{|V(G_i)||V(G_{i+1})|}}$$

**Proof.** First consider the edge  $x_1x_n$ . There are two cases, *n* is even or odd. If n = 2t for some  $t \in \mathbb{N}$ , then, the vertices in the graphs  $G_1, G_2, G_3, \dots, G_t$  are closer to  $x_1$  than  $x_n$ , and the rest are closer to  $x_n$  than  $x_1$ . So,

$$\sqrt{\frac{n_{x_1}(x_1x_n,G) + n_{x_n}(x_1x_n,G) - 2}{n_{x_1}(x_1x_n,G)n_{x_n}(x_1x_n,G)}} = \sqrt{\frac{\sum_{i=1}^t |V(G_i)| + \sum_{i=1}^t |V(G_{t+i})| - 2}{\sum_{i=1}^t |V(G_i)| \sum_{i=1}^t |V(G_{t+i})|}}$$
$$= \sqrt{\frac{|V(G)| - 2}{\sum_{i=1}^t |V(G_i)| \sum_{i=1}^t |V(G_{t+i})|}}$$
$$< \sqrt{\frac{|V(G)| - 2}{|V(G_1)| |V(G_{2t})|}}$$
$$= \sqrt{\frac{|V(G)| - 2}{|V(G_1)| |V(G_n)|}}.$$

By the same argument, for every  $x_i x_{i+1}$ ,  $1 \le i \le n - 1$ , we have:

$$\sqrt{\frac{n_{x_{i}}(x_{i}x_{i+1},G)+n_{x_{i+1}}(x_{i}x_{i+1},G)-2}{n_{x_{i}}(x_{i}x_{i+1},G)n_{x_{i+1}}(x_{i}x_{i+1},G)}} < \sqrt{\frac{|V(G)|-2}{|V(G_{i})||V(G_{i+1})|}}.$$

If n = 2t - 1 for some  $t \in \mathbb{N}$ , then, the vertices in the graphs  $G_1, G_2, G_3, \dots, G_{t-1}$  are closer to  $x_1$  than  $x_n$ , and the vertices in the graphs  $G_{t+1}, G_{t+2}, G_{t+3}, \dots, G_n$  are closer to  $x_n$  than  $x_1$ . The vertices in the graph  $G_t$  have the same distance to  $x_1$  and  $x_n$ . So

$$\frac{\sqrt{\frac{n_{x_1}(x_1x_n,G) + n_{x_n}(x_1x_n,G) - 2}{n_{x_1}(x_1x_n,G)n_{x_n}(x_1x_n,G)}}}{\sqrt{\frac{\sum_{i=1}^{t-1}|V(G_i)| + \sum_{i=1}^{t-1}|V(G_{t+i})| - 2}{\sum_{i=1}^{t-1}|V(G_i)| \sum_{i=1}^{t-1}|V(G_{t+i})|}}} = \sqrt{\frac{|V(G)| - |V(G_t)| - 2}{\sum_{i=1}^{t-1}|V(G_i)| \sum_{i=1}^{t-1}|V(G_{t+i})|}}}.$$

Therefore,

$$\begin{split} \sqrt{\frac{n_{x_1}(x_1x_n,G) + n_{x_n}(x_1x_n,G) - 2}{n_{x_1}(x_1x_n,G)n_{x_n}(x_1x_n,G)}} &< \sqrt{\frac{|V(G)| - |V(G_t)| - 2}{|V(G_1)||V(G_{2t-1})|}} \\ &= \sqrt{\frac{|V(G)| - |V(G_t)| - 2}{|V(G_1)||V(G_n)|}} \\ &< \sqrt{\frac{|V(G)| - 2}{|V(G_1)||V(G_n)|}}. \end{split}$$

By the same argument, for every  $x_i x_{i+1}$ ,  $1 \le i \le n - 1$ , we have:

$$\sqrt{\frac{n_{x_i}(x_ix_{i+1},G) + n_{x_{i+1}}(x_ix_{i+1},G) - 2}{n_{x_i}(x_ix_{i+1},G)n_{x_{i+1}}(x_ix_{i+1},G)}} < \sqrt{\frac{|V(G)| - 2}{|V(G_i)||V(G_{i+1})|}}.$$

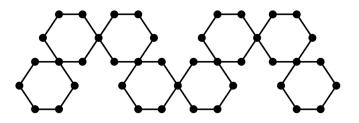
Now by the definition of Graovac-Ghorbani index and similar argument like the proof of the Theorem 2.7, we have the result.

# **3.** CHEMICAL APPLICATIONS

In this section, we apply our previous results in order to obtain the atom-bond connectivity and Graovac-Ghorbani indices of families of graphs that are of importance in chemistry.

### 3.1 Spiro-Chains

Spiro-chains are defined in [6]. Making use of the concept of chain of graphs, a spiro-chain can be defined as a chain of cycles. We denote by  $S_{q,h,k}$  the chain of k cycles  $C_q$  in which the distance between two consecutive contact vertices is h, see  $S_{6,2,8}$  in Figure 7.



**Figure 7**: The graph  $S_{6,2,8}$ .

**Theorem 3.1.** For the graph  $S_{q,h,k}$   $(h \ge 2)$ , we have  $ABC(S_{q,h,k}) = \frac{qk}{\sqrt{2}}$ .

**Proof.** There are 4(k - 1) edges with endpoints of degree 2 and 4. Also there are qk - 4(k - 1) edges with endpoints of degree 2. Therefore

$$ABC(S_{q,h,k}) = 4(k-1)\sqrt{\frac{2+4-2}{2(4)}} + (qk-4(k-1))\sqrt{\frac{2+2-2}{2(2)}},$$

and we have the result.

**Theorem 3.2.** For the graph  $S_{q,1,k}$ , we have  $ABC(S_{q,1,k}) = \frac{qk-k+2}{\sqrt{2}} + \frac{(k-2)\sqrt{6}}{4}$ .

**Proof.** There are k - 2 edges with endpoints of degree 4. Also there are 2k edges with endpoints of degree 4 and 2, and there are qk - 3k + 2 edges with endpoints of degree 2. Therefore by the definition of the atom-bond connectivity, we have the result.

**Theorem 3.3.** Let  $T_n$  be the chain triangular graph of order n. Then,

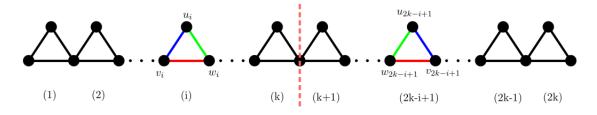
i. for every  $n \ge 2$ , and  $k \ge 1$ , if n = 2k, we have:

$$ABC_{GG}(T_n) = 2\sum_{i=1}^k \left( \sqrt{\frac{2i-2}{2i-1}} + \sqrt{\frac{4k-2i}{4k-2i+1}} + \sqrt{\frac{4k-2}{(4k-2i+1)(2i-1)}} \right),$$

and if n = 2k + 1, then we have:

$$\begin{aligned} ABC_{GG}(T_n) &= 2\sum_{i=1}^k \left( \sqrt{\frac{2i-2}{2i-1}} + \sqrt{\frac{4k-2i+2}{4k-2i+3}} + \sqrt{\frac{4k}{(4k-2i+3)(2i-1)}} \right) \\ &+ 2\sqrt{\frac{2k}{2k+1}} + \frac{2\sqrt{k}}{2k+1}. \end{aligned}$$

ii. for every  $n \ge 2$ ,  $ABC(T_n) = \frac{2n+2}{\sqrt{2}} + \frac{(n-2)\sqrt{6}}{4}$ .



**Figure 8**: Chain triangular cactus  $T_{2k}$ .

**Proof.** (i) We consider the following cases:

**Case 1.** Suppose that *n* is even, and n = 2k for some  $k \in \mathbb{N}$ . Consider the  $T_{2k}$  as shown in Figure 8. One can easily check that whatever happens to computation of Graovac-Ghorbani index related to the edge  $u_i v_i$  in the (*i*)-th triangle in  $T_{2k}$ , is the same as computation of Graovac-Ghorbani index related to the edge  $u_{2k-i+1}v_{2k-i+1}$  in the (2k - i + 1)-th triangle. The same goes for  $w_i v_i$  and  $w_{2k-i+1}v_{2k-i+1}$ , and also for  $w_i u_i$  and  $w_{2k-i+1}u_{2k-i+1}$ . So for computing Graovac-Ghorbani index, it suffices to compute the  $\sqrt{\frac{n_u(uv,T_{2k})+n_v(uv,T_{2k})-2}{n_u(uv,T_{2k})n_v(uv,T_{2k})}}}$  for every  $uv \in E(T_{2k})$  in the first *k* triangles and then multiple that by 2. So from now, we only consider the first *k* triangles.

Consider the blue edge  $u_i v_i$  in the (*i*)-th triangle. There are 2(i-1) + 1 vertices which are closer to  $v_i$  than  $u_i$ , and there is one vertex closer to  $u_i$  than

$$v_i$$
. So,  $\sqrt{\frac{n_{u_i}(u_iv_i,T_{2k}) + n_{v_i}(u_iv_i,T_{2k}) - 2}{n_{u_i}(u_iv_i,T_{2k})n_{v_i}(u_iv_i,T_{2k})}} = \sqrt{\frac{2i-2}{2i-1}}.$ 

Now consider the green edge  $u_i w_i$  in the (*i*)-th triangle. There are 2(2k - i) + 1 vertices which are closer to  $w_i$  than  $u_i$ , and there is one vertex closer to  $u_i$ 

than 
$$w_i$$
. So,  $\sqrt{\frac{n_{u_i}(u_iw_i, T_{2k}) + n_{w_i}(u_iw_i, T_{2k}) - 2}{n_{u_i}(u_iw_i, T_{2k})n_{w_i}(u_iw_i, T_{2k})}} = \sqrt{\frac{4k - 2i}{4k - 2i + 1}}.$ 

Finally, consider the red edge  $v_i w_i$  in the (*i*)-th triangle. There are 2(2k - i) + 1 vertices which are closer to  $w_i$  than  $v_i$ , and there are 2(i - 1) + 1 vertices closer to  $v_i$  than  $w_i$ . So,  $\sqrt{\frac{n_{v_i}(v_i w_i, T_{2k}) + n_{w_i}(v_i w_i, T_{2k}) - 2}{n_{v_i}(v_i w_i, T_{2k})n_{w_i}(v_i w_i, T_{2k})}} = \sqrt{\frac{4k-2}{(4k-2i+1)(2i-1)}}$ .

Since we have k edges like blue one, k edges like green one and k edges like red one, by our argument, we have:

$$ABC_{GG}(T_{2k}) = 2\sum_{i=1}^{k} \left( \sqrt{\frac{2i-2}{2i-1}} + \sqrt{\frac{4k-2i}{4k-2i+1}} + \sqrt{\frac{4k-2}{(4k-2i+1)(2i-1)}} \right).$$

**Case 2.** Suppose that *n* is odd and n = 2k + 1 for some  $k \in \mathbb{N}$ . Now consider the  $T_{2k+1}$  as shown in Figure 9. One can easily check that whatever happens to computation of Graovac-Ghorbani index related to the edge  $u_i v_i$  in the (*i*)-th triangle in  $T_{2k+1}$ , is the same as computation of Graovac-Ghorbani index related to the edge  $u_{2k-i+2}v_{2k-i+2}$  in the (2k - i + 2)-th triangle. The same goes for  $w_i v_i$  and  $w_{2k-i+2}v_{2k-i+2}$ , and also for  $w_i u_i$  and  $w_{2k-i+2}u_{2k-i+2}$ . So for computing Graovac-Ghorbani index, it suffices to compute  $\sqrt{\frac{n_u(uv,T_{2k+1})+n_v(uv,T_{2k+1})-2}{n_u(uv,T_{2k+1})n_v(uv,T_{2k+1})-2}}}$  for every edge  $uv \in E(T_{2k+1})$  in the first *k* triangles and then multiple that by 2 and add it to  $\sum_{uv \in A} \sqrt{\frac{n_u(uv,T_{2k+1})+n_v(uv,T_{2k+1})-2}{n_u(uv,T_{2k+1})n_v(uv,T_{2k+1})}}}$ , where  $A = \{ab, bc, ac\}$ . So from now, we only consider the first *k* triangles and the middle one.

Consider the blue edge  $u_i v_i$  in the (*i*)-th triangle. There are 2(i-1) + 1 vertices which are closer to  $v_i$  than  $u_i$ , and there is one vertex closer to  $u_i$  than

$$v_i$$
. So,  $\sqrt{\frac{n_{u_i}(u_iv_i, T_{2k+1}) + n_{v_i}(u_iv_i, T_{2k+1}) - 2}{n_{u_i}(u_iv_i, T_{2k+1})n_{v_i}(u_iv_i, T_{2k+1})}} = \sqrt{\frac{2i-2}{2i-1}}.$ 

Now consider the green edge  $u_i w_i$  in the (*i*)-th triangle. There are 4k - 2i + 3 vertices which are closer to  $w_i$  than  $u_i$ , and there is one vertex closer to  $u_i$  than  $w_i$ . So,  $\sqrt{\frac{n_{u_i}(u_i w_i, T_{2k+1}) + n_{w_i}(u_i w_i, T_{2k+1}) - 2}{2k + 2}} = \sqrt{\frac{4k - 2i + 2}{2k + 2}}$ 

than 
$$w_i$$
. So,  $\sqrt{\frac{n_{i}}{n_{u_i}(u_iw_i,T_{2k+1})n_{w_i}(u_iw_i,T_{2k+1})}} = \sqrt{\frac{1}{4k-2i+3}}$ .

Next consider the red edge  $v_i w_i$  in the (*i*)-th triangle. There are 2(2k - i + 1) + 1 vertices which are closer to  $w_i$  than  $v_i$ , and there are 2(i - 1) + 1 vertices closer to  $v_i$  than  $w_i$ . So,  $\sqrt{\frac{n_{v_i}(v_i w_i, T_{2k+1}) + n_{w_i}(v_i w_i, T_{2k+1}) - 2}{n_{v_i}(v_i w_i, T_{2k+1}) n_{w_i}(v_i w_i, T_{2k+1})}} = \sqrt{\frac{4k}{(4k - 2i + 3)(2i - 1)}}.$ 

Finally, consider the middle triangle. For the edge ab, there are 2k + 1 vertices which are closer to b than a, and there is one vertex closer to a than b. Also for the edge ac, there are 2k + 1 vertices which are closer to c than a, and there is

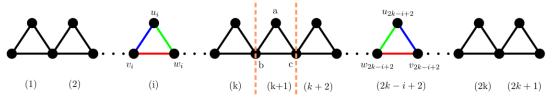
one vertex closer to *a* than *c* and for the edge *bc*, there are 2k + 1 vertices which are closer to *b* than *c*, and there are 2k + 1 vertices closer to *c* than *b*. Hence,  $\sum_{uv \in A} \sqrt{\frac{n_u(uv, T_{2k+1}) + n_v(uv, T_{2k+1}) - 2}{n_u(uv, T_{2k+1})n_v(uv, T_{2k+1})}} = 2\sqrt{\frac{2k}{2k+1}} + \frac{\sqrt{4k}}{2k+1}$ , where  $A = \{ab, bc, ac\}$ .

Since we have k edges like blue one, k edges like green one and k edges like red one, by our argument, we have:

$$ABC_{GG}(T_{2k+1}) = 2\sum_{i=1}^{k} \left( \sqrt{\frac{2i-2}{2i-1}} + \sqrt{\frac{4k-2i+2}{4k-2i+3}} + \sqrt{\frac{4k}{(4k-2i+3)(2i-1)}} \right) + 2\sqrt{\frac{2k}{2k+1}} + \frac{2\sqrt{k}}{2k+1}.$$

Therefore, we have the result.

(ii) It follows from Theorem 3.2.

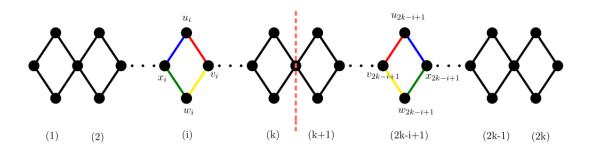


**Figure 9**: Chain triangular cactus  $T_{2k+1}$ .

**Theorem 3.4.** Let  $Q_n$  be the para-chain square cactus graph of order n. Then,

i. for every  $n \ge 1$ , and  $k \ge 1$ , we have:  $ABC_{GG}(Q_n) = \begin{cases} 8\sum_{i=1}^k \sqrt{\frac{6k-1}{(6k-3i+2)(3i-1)}} & \text{if } n = 2k, \\ 8\left(\sum_{i=1}^k \sqrt{\frac{6k+2}{(6k-3i+5)(3i-1)}}\right) + \frac{4\sqrt{6k+2}}{3k+2} & \text{if } n = 2k+1. \end{cases}$ 

ii. for every  $n \ge 2$ ,  $ABC(Q_n) = 2n\sqrt{2}$ .



#### **Figure 10**: Para-chain square cactus $Q_{2k}$ .

#### **Proof.** (i) We consider the following cases:

**Case 1.** Suppose that *n* is even and n = 2k for some  $k \in \mathbb{N}$ . Now consider the  $Q_{2k}$  as shown in Figure 10. One can easily check that whatever happens to computation of Graovac-Ghorbani index related to the edge  $u_i v_i$  in the (*i*)-th rhombus in  $Q_{2k}$ , is the same as computation of Graovac-Ghorbani index related to the edge  $u_{2k-i+1}v_{2k-i+1}$  in the (2k - i + 1)-th rhombus. The same goes for  $w_i v_i$  and  $w_{2k-i+1}v_{2k-i+1}$ , for  $w_i x_i$  and  $w_{2k-i+1}x_{2k-i+1}$ , and also for  $x_i u_i$  and  $x_{2k-i+1}u_{2k-i+1}$ . So for computing Graovac-Ghorbani index, it suffices to compute the  $\sqrt{\frac{n_u(uv,Q_{2k})+n_v(uv,Q_{2k})-2}{n_u(uv,Q_{2k})n_v(uv,Q_{2k})}}$  for every  $uv \in E(Q_{2k})$  in the first *k* rhombus and then multiple that by 2. So from now, we only consider the first *k* rhombus.

Consider the red edge  $u_i v_i$  in the (*i*)-th rhombus. There are 3k + 3(k - i) + 2 vertices which are closer to  $v_i$  than  $u_i$ , and there are 3i - 1 vertices closer to  $u_i$  than  $v_i$ . So,  $\sqrt{\frac{n_{u_i}(u_i v_i, Q_{2k}) + n_{v_i}(u_i v_i, Q_{2k}) - 2}{n_{u_i}(u_i v_i, Q_{2k}) n_{v_i}(u_i v_i, Q_{2k})}} = \sqrt{\frac{6k - 1}{(6k - 3i + 2)(3i - 1)}}.$ 

One can easily check that the edges  $w_i v_i$ ,  $w_i x_i$  and  $x_i u_i$  have the same attitude as  $u_i v_i$ . Since we have k edges like blue one, k edges like green one, k edges like yellow one and k edges like red one, then by our argument, we have  $ABC_{GG}(Q_{2k}) = 2\left(4\sum_{i=1}^k \sqrt{\frac{6k-1}{(6k-3i+2)(3i-1)}}\right)$ .

**Case 2.** Suppose that *n* is odd and n = 2k + 1 for some  $k \in \mathbb{N}$ . Now consider the  $Q_{2k+1}$  as shown in Figure 11. One can easily check that whatever happens to computation of Graovac-Ghorbani index related to the edge  $u_i v_i$  in the (*i*)-th rhombus in  $Q_{2k+1}$ , is the same as computation of Graovac-Ghorbani index related to the edge  $u_{2k-i+2}v_{2k-i+2}$  in the (2k - i + 2)-th rhombus. The same goes for  $w_i v_i$  and  $w_{2k-i+2}v_{2k-i+2}$ , for  $w_i x_i$  and  $w_{2k-i+2}x_{2k-i+2}$ , and also for  $x_i u_i$  and  $x_{2k-i+2}u_{2k-i+2}$ . So for computing Graovac-Ghorbani index, it suffices to compute the  $\sqrt{\frac{n_u(uv,Q_{2k+1})+n_v(uv,Q_{2k+1})-2}{n_u(uv,Q_{2k+1})n_v(uv,Q_{2k+1})}}$  for every  $uv \in E(Q_{2k+1})$  in the first *k* rhombus and then multiple that by 2 and add it to  $\sum_{uv \in A} \sqrt{\frac{n_u(uv,Q_{2k+1})+n_v(uv,Q_{2k+1})-2}{n_u(uv,Q_{2k+1})n_v(uv,Q_{2k+1})}}$ , where  $A = \{ab, bc, cd, da\}$ . So from now, we only consider the first k + 1 rhombus.

Consider the red edge  $u_i v_i$  in the (*i*)-th rhombus. There are 3(k+1) + 3(k-i) + 2 vertices which are closer to  $v_i$  than  $u_i$ , and there are 3i - 1 vertices closer to  $u_i$  than  $v_i$ . So,  $\sqrt{\frac{n_{u_i}(u_i v_i, Q_{2k+1}) + n_{v_i}(u_i v_i, Q_{2k+1}) - 2}{n_{u_i}(u_i v_i, Q_{2k+1}) n_{v_i}(u_i v_i, Q_{2k+1})}} = \sqrt{\frac{6k+2}{(6k-3i+5)(3i-1)}}.$ 

One can easily check that the edges  $w_i v_i$ ,  $w_i x_i$  and  $x_i u_i$  have the same attitude as  $u_i v_i$ .

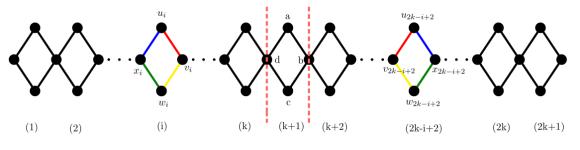
Now consider the middle rhombus. For the edge ab, there are 3k + 2vertices which are closer to b than a, and there are 3k + 2 vertices closer to a than cd edges bc, *b*. the and da have the same attitude as  $\sum_{uv \in A} \sqrt{\frac{n_u(uv, Q_{2k+1}) + n_v(uv, Q_{2k+1}) - 2}{n_u(uv, Q_{2k+1})n_v(uv, Q_{2k+1})}} = \frac{4\sqrt{6k+2}}{3k+2},$ ab. Hence, where  $A = \{ab, bc, cd, da\}.$ 

Since we have k edges like blue one, k edges like green one, k edges like yellow one and k edges like red one, then by our argument, we have:

$$ABC_{GG}(Q_{2k+1}) = 2\left(4\sum_{i=1}^{k}\sqrt{\frac{6k+2}{(6k-3i+5)(3i-1)}}\right) + \frac{4\sqrt{6k+2}}{3k+2}$$

Therefore, we have the result.

(ii) It follows from Theorem 3.1.



**Figure 11**: Para-chain square cactus  $Q_{2k+1}$ .

**Theorem 3.5.** Let  $O_n$  be the para-chain square cactus graph of order n. Then, i. for every  $n \ge 2$ , and  $k \ge 1$ , if n = 2k, we have:

$$ABC_{GG}(O_n) = 2k\sqrt{2} + 4\left(\sum_{i=1}^k \sqrt{\frac{6k-1}{(6k-3i+2)(3i-1)}}\right),$$

and if n = 2k + 1, then we have:

$$ABC_{GG}(O_n) = (2k+1)\sqrt{2} + \frac{2\sqrt{6k+2}}{3k+2} + 4\left(\sum_{i=1}^k \sqrt{\frac{6k+2}{(6k-3i+5)(3i-1)}}\right)$$
  
.ii for every  $n \ge 2$ ,  $ABC(O_n) = \frac{3n+2}{\sqrt{2}} + \frac{(n-2)\sqrt{6}}{4}$ .

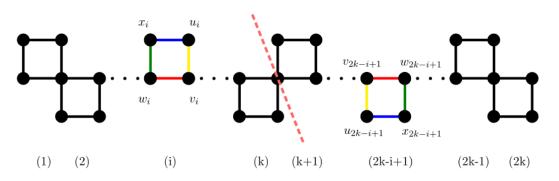


Figure 12: Para-chain square cactus  $O_{2k}$ .

**Proof.** (i)We consider the following cases:

**Case 1.** Suppose that *n* is even and n = 2k for some  $k \in \mathbb{N}$ . Now consider the  $O_{2k}$  as shown in Figure 12. One can easily check that whatever happens to computation of Graovac-Ghorbani index related to the edge  $u_i v_i$  in the (*i*)-th square in  $O_{2k}$ , is the same as computation of Graovac-Ghorbani index related to the edge  $u_{2k-i+1}v_{2k-i+1}$  in the (2k - i + 1)-th square. The same goes for  $w_i v_i$  and  $w_{2k-i+1}v_{2k-i+1}$ , for  $w_i x_i$  and  $w_{2k-i+1}x_{2k-i+1}$ , and also for  $x_i u_i$  and  $x_{2k-i+1}u_{2k-i+1}$ . So for computing Graovac-Ghorbani index, it suffices to compute the  $\sqrt{\frac{n_u(uv,O_{2k})+n_v(uv,O_{2k})-2}{n_u(uv,O_{2k})n_v(uv,O_{2k})}}$  for every  $uv \in E(O_{2k})$  in the first *k* squares and then multiple that by 2. So from now, we only consider the first *k* squares.

Consider the yellow edge  $u_i v_i$  in the (*i*)-th square. There are 3(2k) - 1 vertices which are closer to  $v_i$  than  $u_i$ , and there are 2 vertices closer to  $u_i$  than  $v_i$  which is  $x_i$ . So,  $\sqrt{\frac{n_{u_i}(u_i v_i, O_{2k}) + n_{v_i}(u_i v_i, O_{2k}) - 2}{n_{u_i}(u_i v_i, O_{2k})n_{v_i}(u_i v_i, O_{2k})}} = \frac{\sqrt{2}}{2}$ . By the same argument, the same happens to the edge  $x_i w_i$ .

Now consider the blue edge  $u_i x_i$  in the (*i*)-th square. There are 3i - 1 vertices which are closer to  $x_i$  than  $u_i$ , and there are 3k + 3(k - i) + 2 vertices closer to  $u_i$  than  $x_i$ . So,  $\sqrt{\frac{n_{u_i}(u_i x_i, O_{2k}) + n_{x_i}(u_i x_i, O_{2k}) - 2}{n_{u_i}(u_i x_i, O_{2k})n_{x_i}(u_i x_i, O_{2k})}} = \sqrt{\frac{6k-1}{(6k-3i+2)(3i-1)}}$ . By the same argument, the same happens to the edge  $v_i w_i$ .

Since we have k edges like blue one, k edges like green one, k edges like yellow one and k edges like red one, then by our argument, we have:

$$ABC_{GG}(O_{2k}) = 2\left(2\sum_{i=1}^{k} \frac{\sqrt{2}}{2} + 2\sum_{i=1}^{k} \sqrt{\frac{6k-1}{(6k-3i+2)(3i-1)}}\right)$$

**Case 2.** Suppose that *n* is odd and n = 2k + 1 for some  $k \in \mathbb{N}$ . Now consider the  $O_{2k+1}$  as shown in Figure 13. One can easily check that whatever happens to computation of Graovac-Ghorbani index related to the edge  $u_i v_i$  in the (*i*)-th square in  $O_{2k+1}$ , is the same as computation of Graovac-Ghorbani index related to

the edge  $u_{2k-i+2}v_{2k-i+2}$  in the (2k-i+2)-th square. The same goes for  $w_iv_i$  and  $w_{2k-i+2}v_{2k-i+2}$ , for  $w_ix_i$  and  $w_{2k-i+2}x_{2k-i+2}$ , and also for  $x_iu_i$  and  $x_{2k-i+2}u_{2k-i+2}$ . So for computing Graovac-Ghorbani index, it suffices to compute the  $\sqrt{\frac{n_u(uv,O_{2k+1})+n_v(uv,O_{2k+1})-2}{n_u(uv,O_{2k+1})n_v(uv,O_{2k+1})}}$  for every  $uv \in E(O_{2k+1})$  in the first k squares and then multiple that by 2 and add it to  $\sum_{uv \in A} \sqrt{\frac{n_u(uv, O_{2k+1}) + n_v(uv, O_{2k+1}) - 2}{n_u(uv, O_{2k+1})n_v(uv, O_{2k+1})}}$ , where

 $A = \{ab, bc, cd, da\}$ . So from now, we only consider the first k + 1 squares.

Consider the yellow edge  $u_i v_i$  in the (i)-th square. There are 3(2k + 1) - 1vertices which are closer to  $v_i$  than  $u_i$ , and there are 2 vertices closer to  $u_i$  than  $v_i$ . So,  $\sqrt{\frac{n_{u_i}(u_iv_i, o_{2k+1}) + n_{v_i}(u_iv_i, o_{2k+1}) - 2}{n_{u_i}(u_iv_i, o_{2k+1})n_{v_i}(u_iv_i, o_{2k+1})}} = \frac{\sqrt{2}}{2}$ . By the same argument, the same happens to the edge  $x_i w_i$ .

Now consider the blue edge  $u_i x_i$  in the (i)-th square. There are 3i - 1vertices which are closer to  $x_i$  than  $u_i$ , and there are 3(k+1) + 3closer to  $u_i$  than  $x_i$ . So,  $\sqrt{\frac{n_{u_i}(u_i x_i, o_{2k+1}) + n_{x_i}(u_i x_i, o_{2k+1}) - 2}{n_{u_i}(u_i x_i, o_{2k+1})n_{x_i}(u_i x_i, o_{2k+1})}} =$ vertices  $\sqrt{\frac{6k+2}{(6k-3i+5)(3i-1)}}$ . By the same argument, the same happens to the edge  $v_i w_i$ .

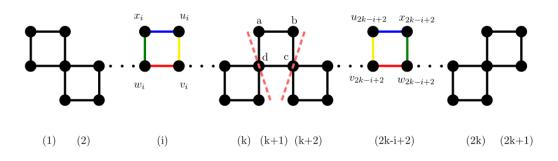
Now consider the middle square. For the edge ab, there are 3k + 2 vertices which are closer to b than a, and there are 3k + 2 vertices closer to a than b. The edge cd has the same attitude as ab. But for the edge ad, there are 3(2k + 1) - 1vertices which are closer to d than a, and there are 2 vertices closer to a than d, and edge bc the same attitude as Hence,  $\sum_{uv \in A} \sqrt{\frac{n_u(uv, O_{2k+1}) + n_v(uv, O_{2k+1}) - 2}{n_u(uv, O_{2k+1})n_v(uv, O_{2k+1})}} = \frac{2\sqrt{6k+2}}{3k+2} + \sqrt{2},$ where A =ad.  $\{ab, bc, cd, da\}.$ 

Since we have k edges like blue one, k edges like green one, k edges like yellow one and k edges like red one, then by our argument, we have:

$$ABC_{GG}(O_{2k+1}) = 2\left(2\sum_{i=1}^{k} \frac{\sqrt{2}}{2} + 2\sum_{i=1}^{k} \sqrt{\frac{6k+2}{(6k-3i+5)(3i-1)}}\right) + \frac{2\sqrt{6k+2}}{3k+2} + \sqrt{2}.$$

Therefore, we have the result.

(ii) It follows from Theorem 3.2.



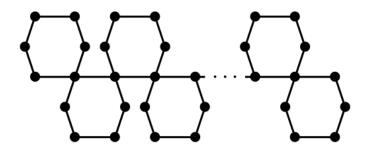
**Figure 13**: Para-chain square cactus  $O_{2k+1}$ .

**Theorem 3.6.** Let  $O_n^h$  be the Ortho-chain graph of order n (see Figure 14). Then,

i. for every 
$$n \ge 2$$
, and  $k \ge 1$ , if  $n = 2k$ , then we have:  
 $ABC_{GG}(O_n^h) = 4\left(\sum_{i=1}^k \sqrt{\frac{10k-1}{(10k-5i+3)(5i-2)}}\right) + 8k\sqrt{\frac{10k-1}{30k-6}}$ ,  
and if  $n = 2k + 1$ , we have:  
 $ABC_{GG}(O_n^h) = 4\left(\sum_{i=1}^k \sqrt{\frac{10k+4}{(10k-5i+8)(5i-2)}}\right) + (8k + 4)\sqrt{\frac{10k+4}{30k+9}} + \frac{2\sqrt{10k+4}}{5k+3}$ 

ii. for every 
$$n \ge 2$$
,  $ABC(O_n^h) = \frac{5n+2}{\sqrt{2}} + \frac{(n-2)\sqrt{6}}{4}$ .

**Proof.** (i) It is similar to the proof of Theorem 3.5. (ii) It follows from Theorem 3.2.

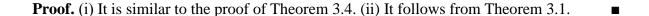


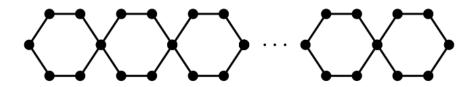
**Figure 14**: Ortho-chain graph  $O_n^h$ .

**Theorem 3.7.** Let  $L_n$  be the para-chain hexagonal graph of order n (see Figure 15). Then, i. for every  $n \ge 1$ , and  $k \ge 1$ , we have:

$$ABC_{GG}(L_n) = \begin{cases} 12\sum_{i=1}^k \sqrt{\frac{10k-1}{(10k-5i+3)(5i-2)}} & \text{if } n = 2k, \\ 12\left(\sum_{i=1}^k \sqrt{\frac{10k+4}{(10k-5i+8)(5i-2)}}\right) + \frac{6\sqrt{10k+4}}{5k+3} & \text{if } n = 2k+1. \end{cases}$$
  
for every  $n \ge 2$ ,  $ABC(L_n) = 3n\sqrt{2}$ .

ii.





**Figure 15**: Para-chain hexagonal graph  $L_n$ .

**Theorem 3.8.** Let  $M_n$  be the Meta-chain hexagonal of order n (see Figure 16). Then,

i. for every  $n \ge 2$ , and  $k \ge 1$ , if n = 2k, we have:

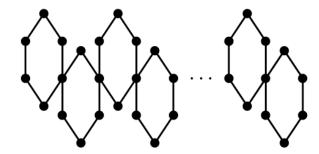
$$ABC_{GG}(M_n) = 8\left(\sum_{i=1}^k \sqrt{\frac{10k-1}{(10k-5i+3)(5i-2)}}\right) + 4k\sqrt{\frac{10k-1}{30k-6}}$$

and if n = 2k + 1, then we have:

$$ABC_{GG}(M_n) = 8\left(\sum_{i=1}^k \sqrt{\frac{10k+4}{(10k-5i+8)(5i-2)}}\right) + (2k+2)\sqrt{\frac{10k+4}{30k+9}} + \frac{4\sqrt{10k+4}}{5k+3}$$
  
ii. for every  $n \ge 2$ ,  $ABC(M_n) = 3n\sqrt{2}$ .

**Proof.** (i) It is similar to the proof of Theorem 3.5. (ii) It follows from Theorem 3.1.

**Corollary 3.9.** *Meta-chain hexagonal cactus graphs and para-chain hexagonal cactus graphs of the same order, have the same atom-bond connectivity index. But they do not have the same Graovac-Ghorbani index.* 



**Figure 16**: Meta-chain hexagonal graph  $M_n$ .

# **3.2 POLYPHENYLENES**

Similar to the above definition of the spiro-chain  $S_{q,h,k}$ , we can define the graph  $L_{q,h,k}$  as the link of k cycles  $C_q$  in which the distance between the two contact vertices in the same cycle is h, see  $L_{6,2,4}$  in Figure 17.

**Theorem 3.10.** For the graph  $L_{q,h,k}$ , when  $h \ge 2$ , we have  $ABC(L_{q,h,k}) = \frac{2(k-1)}{3} + \frac{qk}{\sqrt{2}}$ .

**Proof.** There are k-1 edges with endpoints of degree 3. Also there are 4(k-1) edges with endpoints of degree 3 and 2, and there are qk - 4(k-1) edges with endpoints of degree 2. Therefore

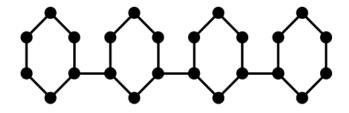
$$ABC(L_{q,h,k}) = (k-1)\sqrt{\frac{3+3-2}{3(3)}} + 4(k-1)\sqrt{\frac{3+2-2}{3(2)}} + (qk-4(k-1))\sqrt{\frac{2+2-2}{2(2)}},$$
  
have the result.

and we

**Theorem 3.11.** For the graph  $L_{a,1,k}$ , we have:

$$ABC(L_{q,1,k}) = \frac{4k-6}{3} + \frac{qk-k+2}{\sqrt{2}}.$$

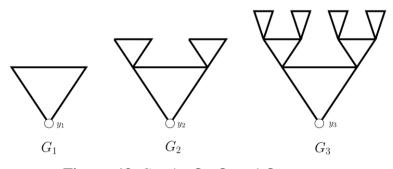
**Proof.** There are 2k - 3 edges with endpoints of degree 3. Also there are 2k edges with endpoints of degree 3 and 2, and there are qk - 3k + 2 edges with endpoints of degree 2. Therefore, by the definition of the atom-bond connectivity index, we have the result.



**Figure 17**: The graph  $L_{6,2,4}$ .

## **3.3 TRIANGULANES**

We intend to derive the atom-bond connectivity of the triangulane  $T_k$  defined pictorially in [19]. We define  $T_k$  recursively in a manner that will be useful in our approach. First we define recursively an auxiliary family of triangulanes  $G_k$  ( $k \ge 1$ ). Let  $G_1$  be a triangle and denote one of its vertices by  $y_1$ . We define  $G_k$   $(k \ge 2)$  as the circuit of the graphs  $G_{k-1}$ ,  $G_{k-1}$ , and  $K_1$  and denote by  $y_k$  the vertex where  $K_1$  has been placed. The graphs  $G_1$ ,  $G_2$  and  $G_3$  are shown in Figure 18.



**Figure 18**: Graphs  $G_1$ ,  $G_2$  and  $G_3$ .

**Theorem 3.12.** For the graph  $T_k$  (see  $T_3$  in Figure 19), we have:

i. 
$$ABC(T_k) = \frac{9(2^{k-1})\sqrt{2}}{2} + \frac{(9(2^k)-6)\sqrt{6}}{4}.$$
  
ii. 
$$ABC_{GG}(T_n) = 6\sqrt{\frac{2^{n+2}+2^{n-4}}{(2^{n+2}-1)(2^{n-1})}} + \frac{3\sqrt{2^{n+2}-4}}{2^{n+1}-1}$$
$$+ \sum_{i=2}^n 3(2^i) \left(\sqrt{\frac{2^{n+2}-(\sum_{i=2}^{i-2}2^{n-i})+2^{n-i+1}-4}{(2^{n+2}-1+\sum_{i=2}^{i-2}2^{n-i})(2^{n-i+1}-1)}}\right)$$
$$+ \sum_{i=1}^n 3(2^{i-1}) \left(\frac{\sqrt{2^{n-i+2}-4}}{2^{n-i+1}-1}\right).$$

**Figure 19**: Graph *T*<sub>3</sub>.

**Proof.** (i) Since creating such a graph is recursive, then there are  $3 + 3\sum_{n=0}^{k-1} 3(2^n)$  edges with endpoints of degree 4. Also there are  $3(2^k)$  edges with endpoints of degree 4 and 2, and there are  $3(2^{k-1})$  edges with endpoints of degree 2. Therefore, by the definition of the atom-bond connectivity index, and we have the result.

(ii) Consider the graph  $T_n$  in Figure 20. First we consider the edge  $x_0x_1$ . There are  $2^{n+2} - 1$  vertices which are closer to  $x_0$  than  $x_1$ , and there are  $2^n - 1$  vertices closer to  $x_1$  than  $x_0$ . So,  $\sqrt{\frac{n_{x_0}(x_0x_1,T_n)+n_{x_1}(x_0x_1,T_n)-2}{n_{x_0}(x_0x_1,T_n)n_{x_1}(x_0x_1,T_n)}} = \sqrt{\frac{2^{n+2}+2^{n-4}}{(2^{n+2}-1)(2^{n}-1)}}$ . The edge  $ax_0$  has the same attitude as the blue edge  $x_0x_1$ . In total there are 6 edges with this value related to Graovac-Ghorbani index. The number of vertices closer to vertex a is the same as the number of vertices closer to vertex  $x_1$  and are  $2^n - 1$  vertices. So,  $\sqrt{\frac{n_a(ax_1,T_n)+n_{x_1}(ax_1,T_n)-2}{n_a(ax_1,T_n)n_{x_1}(ax_1,T_n)}} = \frac{\sqrt{2^{n+1}-4}}{2^{n}-1}$ , and in total, we have 3 edges like this one.

Now consider the edge  $x_1x_2$ . There are  $2(2^{n+1}-1) + 2^n + 1$  vertices which are closer to  $x_1$  than  $x_2$ , and there are  $2^{n-1} - 1$  vertices closer to  $x_2$  than

 $x_1$ . So,  $\sqrt{\frac{n_{x_0}(x_0x_1,T_n)+n_{x_1}(x_0x_1,T_n)-2}{n_{x_0}(x_0x_1,T_n)n_{x_1}(x_0x_1,T_n)}} = \sqrt{\frac{2^{n+2}+2^{n+2n-1}-4}{(2^{n+2}+2^{n-1}-1)(2^{n-1}-1)}}$ . The edge  $bx_1$  has the same attitude as the red edge  $x_1x_2$ . In total there are 12 edges with this value related to Graovac-Ghorbani index. The number of vertices closer to vertex b is the same as the number of vertices closer to vertex  $x_2$ , and in total, and are  $2^{n-1}-1$  vertices. So,  $\sqrt{\frac{n_b(bx_1,T_n)+n_{x_1}(bx_1,T_n)-2}{n_b(bx_1,T_n)n_{x_1}(bx_1,T_n)}} = \frac{\sqrt{2^{n-4}}}{2^{n-1}-1}$ , and in total, we have 6 edges like this one.

By continuing this process in the *i*-th level (i > 1), we have:

$$\sqrt{\frac{n_{x_{i-1}}(x_{i-1}x_i,T_n) + n_{x_i}(x_{i-1}x_i,T_n) - 2}{n_{x_{i-1}}(x_{i-1}x_i,T_n)n_{x_i}(x_{i-1}x_i,T_n)}} = \sqrt{\frac{2^{n+2} + (\sum_{t=0}^{i-2} 2^{n-t}) + 2^{n-i+1} - 4}{(2^{n+2} - 1 + \sum_{t=0}^{i-2} 2^{n-t})(2^{n-i+1} - 1)}}$$

We have  $3(2^{i})$  edges like this one. The number of vertices closer to vertex  $x_{i}$  is the same as the number of vertices closer to its neighbour in horizontal edge with one endpoint  $x_{i}$  (suppose *l*), and are  $2^{n-i+2} - 1$  vertices. So,  $\sqrt{\frac{n_{l}(lx_{1},T_{n})+n_{x_{1}}(lx_{1},T_{n})-2}{n_{l}(lx_{1},T_{n})n_{x_{1}}(lx_{1},T_{n})}} = \frac{\sqrt{2^{n-i+2}-4}}{2^{n-i+1}-1}$ , and in total, we have  $3(2^{i-1})$  edges like this one.

Finally, the number of vertices closer to vertex  $x_0$  is the same as the number of vertices closer to vertex u, the number of vertices closer to vertex  $x_0$  is the same as the number of vertices closer to vertex v, and the number of vertices closer to vertex v is the same as the number of vertices closer to vertex u, and are  $2^{n+1} - 1$  vertices. So by the definition of the Graovac-Ghorbani index and our argument, we have

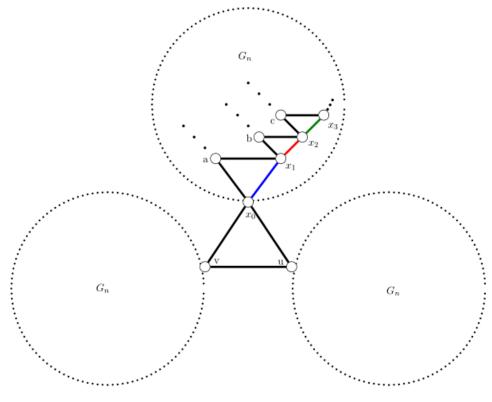
$$\begin{aligned} ABC_{GG}(T_n) &= 6\sqrt{\frac{2^{n+2}+2^{n-4}}{(2^{n+2}-1)(2^{n-1})}} \\ &+ \sum_{i=2}^n 3(2^i) \left(\sqrt{\frac{2^{n+2}+(\sum_{t=0}^{i-2}2^{n-t})+2^{n-i+1}-4}{(2^{n+2}-1+\sum_{t=0}^{i-2}2^{n-t})(2^{n-i+1}-1)}}\right) \\ &+ \left(\sum_{i=1}^n 3(2^{i-1}) \left(\frac{\sqrt{2^{n-i+2}-4}}{2^{n-i+1}-1}\right)\right) + \frac{3\sqrt{2^{n+2}-4}}{2^{n+1}-1}, \end{aligned}$$

and therefore we have the result.

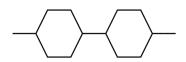
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#### **3.4 NANOSTAR DENDRIMERS**

We want to compute the atom-bond connectivity of the nanostar dendrimer  $D_k$  defined in [19]. First we define recursively an auxiliary family of rooted dendrimers  $G_k$  ( $k \ge 1$ ). We need a fixed graph F defined in Figure 21, we consider one of its endpoint to be the root of F. The graph  $G_1$  is defined in Figure 21, the leaf being its root. Now we define  $G_k$  ( $k \ge 2$ ) the bouquet of the following 3 graphs:  $G_{k-1}, G_{k-1}$ , and F with respect to their roots; the root of  $G_k$  is taken to be its unique leaf (see  $G_2$  and  $G_3$  in Figure 22). Finally, we define  $D_k$  ( $k \ge 1$ ) as the bouquet of 3 copies of  $G_k$  with respect to their roots ( $D_2$  is shown in Figure 23, where the circles represent hexagons).

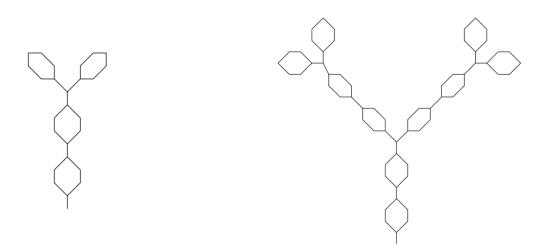


**Figure 20**: Graph  $T_n$ .





**Figure 21**: Graphs F and  $G_1$ , respectively.



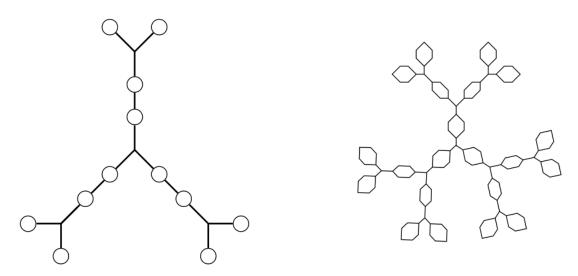
**Figure 22**: Graphs  $G_2$  and  $G_3$ , respectively.

**Theorem 3.13.** For the dendrimer  $D_3[n]$  we have:  $ABC(D_3[n]) = 6(2^n) - 4 + (18(2^n) - 9)\sqrt{2}.$ 

**Proof.** There are  $3 + 9 \sum_{k=0}^{n-1} 2^k$  edges with endpoints of degree 3. Also there are  $6 + 18 \sum_{k=0}^{n-1} 2^k$  edges with endpoints of degree 3 and 2, and there are  $12 + 18 \sum_{k=0}^{n-1} 2^k$  edges with endpoints of degree 2. Therefore

$$\begin{split} ABC(D_3[n]) &= (3+9\sum_{k=0}^{n-1}2^k)\sqrt{\frac{3+3-2}{3(3)}} + (6+18\sum_{k=0}^{n-1}2^k)\sqrt{\frac{3+2-2}{3(2)}} \\ &+ (12+18\sum_{k=0}^{n-1}2^k)\sqrt{\frac{2+2-2}{2(3)}}, \end{split}$$

and we have the result.



**Figure 23**: Nanostar  $D_2$  and  $D_3$ [2], respectively.

**ACKNOWLEDGEMENTS.** The author would like to express his gratitude to the referees for helpful comments. Also, he would like to thank the Research Council of Norway and Department of Informatics, University of Bergen for their support.

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