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# The Second Geometric–Arithmetic Index for Trees and Unicyclic Graphs

# NASRIN DEHGARDI<sup>a</sup>, HAMIDEH ARAM<sup>b</sup> AND ABDOLLAH KHODKAR<sup>•</sup>

<sup>a</sup>Department of Mathematics and Computer Science, Sirjan University of Technology, Sirjan, Iran

<sup>b</sup>Department of Mathematics, Gareziaeddin Center, Khoy Branch, Islamic Azad University, Khoy, Iran

<sup>•</sup>Department of Mathematics, University of West Georgia, Carrollton, GA 30118, USA

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#### ABSTRACT

Let G be a finite and simple graph with edge set E(G). The second geometric-arithmetic Index is defined as

$$GA_2(G) = \sum_{uv \in E(G)} \frac{2\sqrt{n_u n_v}}{n_u + n_v}$$

where  $n_u$  denotes the number of vertices in *G* lying closer to *u* than to v. In this paper we find a sharp upper bound for  $GA_2(T)$ , where *T* is tree, in terms of the order and maximum degree of the tree. We also find a sharp upper bound for  $GA_2(G)$ , where *G* is a unicyclic graph, in terms of the order, maximum degree and girth of *G*. In addition, we characterize the trees and unicyclic graphs which achieve the upper bounds.

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### **1 INTRODUCTION**

Let G be a simple connected graph with vertex set V = V(G) and edge set E = E(G). The integers n = n(G) = |V(G)| and m = m(G) = |E(G)| are the order and the size of the graph G, respectively. We write  $deg_G(v) = deg(v)$  for the degree of a vertex v and  $\Delta = \Delta(G)$  for the maximum degree of G. Let  $u, v \in V(G)$ , then the distance  $d_G(u, v)$  between u and v is defined as the length of a shortest path in G connecting u and v.

<sup>•</sup> Corresponding Author (Email address:akhodkar@westga.edu) DOI: 10.22052/ijmc.2017.81079.1277

In [5], a new class of topological descriptors, based on some properties of the vertices of a graph is presented. These descriptors are named as geometric-arithmetic indices,  $GA_{general}$ , and defined as:

$$GA_{\text{general}}(G) = \sum_{uv \in E(G)} \frac{2\sqrt{Q_u Q_v}}{Q_u + Q_v},$$

where  $Q_u$  is some quantity that in a unique manner can be associated with the vertex u of the graph G. The geometric-arithmetic index GA is defined in [6] as:

$$GA(G) = \sum_{uv \in E(G)} \frac{2\sqrt{deg(u)deg(v)}}{deg(u) + deg(v)}$$

The geometric-arithmetic index is well studied in the literature, see for example [2, 4, 7]. Let uv be an edge of G. Define  $N(u, G) = \{x \in V(G) \mid d_G(u, x) < d_G(u, x)\}$ . In other words, N(u, G) consists of vertices of G which are closer to u than to v. Note that the vertices equidistant to u and v are not included into either N(u, G) or N(v, G). Such vertices exist only if the edge uv belongs to an odd cycle. Hence, in trees,  $n_u + n_v = n$  for all edges of the tree. It is also worth noting that  $u \in N(u, G)$  and  $v \in N(v, G)$ , which implies that  $n_u \ge 1$  and  $n_v \ge 1$ . The second geometric-arithmetic index  $GA_2$  is defined in [5] as:

$$GA_2(G) = \sum_{uv \in E(G)} \frac{2\sqrt{n_u n_v}}{n_u + n_v}$$

where  $n_u = n_u(G) = |N(u, G)|$ . See [1, 3, 8] for more information on this index.

The following statements can be found in [5].

**Theorem A.** The path  $P_n$  is the n-vertex tree with maximum second geometric-arithmetic index.

**Theorem B.** Let  $S_n$  be a star of order n, then  $GA_2(G) = \frac{2(n-1)\sqrt{n-1}}{n}$ .

In this paper we first present some examples. Then we prove that for any tree *T* of order  $n \ge 2$  with maximum degree  $\Delta$ ,

$$GA_2(T) \le \frac{2}{n} \left( (\Delta - 1)\sqrt{n - 1} + \sum_{i=1}^{n - \Delta} \sqrt{i(n - i)} \right)$$

Finally, we prove that for any unicyclic graph *G* of order  $n \ge 3$  with maximum degree  $\Delta \ge 3$  and girth *k*, if *k* is odd, then

$$\begin{aligned} GA_2(G) &\leq \frac{2}{n} \left( (\Delta - 2)\sqrt{n - 1} + \sum_{i=1}^{n-k-\Delta+2} \sqrt{i(n-i)} \right) \\ &+ \frac{2(k-1)}{n-1} \sqrt{\left(\frac{k-1}{2} + \Delta - 2\right)\left(n - \frac{k-1}{2} - \Delta + 1\right)} + \frac{2}{\Delta + k - 3} \sqrt{\frac{k-1}{2}\left(\frac{k-1}{2} + \Delta - 2\right)} \\ &+ \frac{2}{n-\Delta+1} \sqrt{\frac{k-1}{2}\left(n - \frac{k-1}{2} - \Delta + 1\right)}, \end{aligned}$$

and if *k* is even, then

$$GA_{2}(G) \leq \frac{2}{n} \left( (\Delta - 2)\sqrt{n - 1} + \sum_{i=1}^{n - k - \Delta + 2} \sqrt{i(n - i)} + k \sqrt{(\frac{k}{2} + \Delta - 2)(n - \frac{k}{2} - \Delta + 2)} \right)$$

We also characterize the trees and unicyclic graphs which achieve the upper bounds.

## 2 **EXAMPLES**

Dendrimers are nanostructures that can be precisely designed and manufactured for a wide variety of applications, such as drug delivery, gene delivery and diagnostic tests. In this section we calculate the second geometric-arithmetic index for Dendrimers of types A and B and for Tecto Dendrimers. See Figure 1.



Figure 1: Dendrimers of types A and B and Tecto Dendrimers.

**Example 1.** In Dendrimers D[n] type A, denoted  $D[n]_A$ , there are  $4(2^n - 1) + 1$  vertices and  $4(2^n - 1)$  edges. Let e be an edge between the *i*th and the (i + 1)th layers. Then

$$f_i(e) = \sqrt{(2^{n-i} - 1)(2^{n+2} - 2^{n-i} - 2)} \text{ for } i = 1, 2, \dots, n-1.$$

In addition, there are  $2^{i+2}$  edges between the *i*th and the (i + 1)th layers. Therefore, for  $n \ge 2$ ,

$$\begin{aligned} GA_2(D[n]_A) &= \frac{2}{4(2^n-1)+1} \bigg( 4\sqrt{(2^n-1)(3(2^n-1)+1)} + \sum_{i=1}^{n-1} 2^{i+2} f_i(e) \bigg) \\ &= \frac{8}{4(2^n-1)+1} \bigg( \sqrt{(2^n-1)(3(2^n-1)+1)} + \sum_{i=1}^{n-1} 2^i f_i(e) \bigg). \end{aligned}$$

For examples,

$$GA_2(D[2]_A) = \frac{8}{13} \left(\sqrt{30} + 2\sqrt{12}\right) = 7.63$$
 and  
 $GA_2(D[3]_A) = \frac{8}{29} \left(\sqrt{154} + 2\sqrt{78} + 4\sqrt{28}\right) = 14.13.$ 

**Example 2.** In Dendrimers D[n] type *B*, denoted  $D[n]_{B}$ , there are  $3(2^n - 1) + 1$  vertices and  $3(2^n - 1)$  edges. Let *e* be an edge between the *i*th and the (i + 1)th layers. Then

$$f_i(e) = \sqrt{(2^{n-i}-1)(3(2^n)-2^{n-i}-1)}$$
 for  $i = 1, 2, ..., n-1$ .

In addition, there are  $3(2^i)$  edges between the *i*th and the (i + 1)th layers. Therefore, for  $n \ge 2$ ,

$$GA_2(D[n]_B) = \frac{8}{3(2^n-1)+1} \left( \sqrt{3(2^n-1)(2(2^n-1)+1)} + \sum_{i=1}^{n-1} 3(2^i) f_i(e) \right)$$

For example,

$$GA_2(D[2]_B) = \frac{2}{10} (3\sqrt{21} + 18) = 6.35$$
 and  
 $GA_2(D[3]_B) = \frac{2}{22} (3\sqrt{105} + 6\sqrt{57} + 12\sqrt{21}) = 11.91.$ 

**Example 3.** In Tecto Dendrimers  $D[n]_T$ , there are  $2^{n+2} - 2$  vertices and  $2^{n+2} - 3$  edges. Let *e* be an edge between the *i*th and the (i + 1)th layers. Then

$$f_i(e) = \sqrt{(2^{n-i} - 1)(2^{n+2} - 2^{n-i} - 1)}$$
 for  $i = 1, 2, ..., n - 1$ 

In addition, there are  $2^{i+2}$  edges between the *i*th and the (i + 1)th layers. Therefore, for  $n \ge 2$ ,

$$GA_2(D[n]_T) = \frac{8}{2^{n+2}-2} \Big( 4\sqrt{(2^n-1)(3(2^n)-1)} + \sum_{i=1}^{n-1} 2^{i+2} f_i(e) + 2^{n+1} - 1 \Big).$$

For example,

$$GA_2(D[2]_T) = \frac{2}{14} \left( 4\sqrt{33} + 8\sqrt{13} + 7 \right) = 8.40 \text{ and}$$
  

$$GA_2(D[3]_T) = \frac{2}{30} \left( 4\sqrt{161} + 8\sqrt{81} + 16\sqrt{29} + 15 \right) = 14.93.$$

# 2 AN UPPER BOUND ON THE SECOND GEOMETRIC-ARITHMETIC OF TREES

In this section we present a sharp upper bound for the second geometric-arithmetic index of trees in terms of their order and maximum degree. We also characterize all trees whose the second geometric-arithmetic index achieves the upper bound. A *leaf* of a tree *T* is a vertex of degree 1, a *support vertex* is a vertex adjacent to a leaf and a *strong support vertex* is a support vertex adjacent to at least two leaves. An end-support vertex is a support vertex whose all neighbors with exception at most one are leaves. A *rooted tree* is a tree having a distinguished vertex *v*, called the *root*. Let  $T_{n,\Delta}$  be the set of trees of order *n* and maximum degree  $\Delta$ . Let *T* be a tree of order *n* and let  $f : E(T) \rightarrow Z^+$  is a function defined by  $f(xy) = \sqrt{n_x n_y}$ . Hence  $GA_2(T) = \frac{2}{n} \sum_{uv \in E(G)} f(uv)$ . We start with an easy but useful observation.

**Observation 4.** Let  $x \ge y \ge 1$  and  $n \ge x + y + 2$  be positive integers. Then for every  $1 \le k \le y$ , (x + k)(n - x - k) > (y - k + 1)(n - y + k - 1).

*Proof.* First note that  $(x + k)(n - x - k) - (y - k + 1)(n - y + k - 1) = n(2k + x) - (x + k)^2$ . Since  $n \ge x + y + 2$ , it follows that  $n(2k + x) - (x + k)^2 > 0$ . So the result follows.

**Lemma 5.** Let *T* be a tree of order *n* with maximum degree  $\Delta$  and *v* be a vertex of maximum degree. If *T* has a vertex of degree at least three different from *v*, then there is a tree  $T' \in T_{n,\Delta}$  such that  $GA_2(T) < GA_2(T')$ .

*Proof.* Let T be the rooted tree at v. Let  $u \neq v$  be a vertex of degree  $deg(u) = k \geq 3$  such that d(u, v) is as large as possible and let  $N(u) = \{u_1, u_2, \dots, u_{k-2}, u_{k-1}, u_k\}$ . Now we distinguish three cases.

Case 1. *u* is an end-support vertex.

We may assume that  $u_k$  is the parent of u. Let  $S = \{uu_1, uu_2, \dots, uu_{k-2}, uu_{k-1}\}$  and let T' be the tree obtained by attaching the path  $uu_1u_2 \dots u_{k-2}u_{k-1}$  to  $T - \{u_1, u_2, \dots, u_{k-1}\}$ . Suppose that  $S' = \{uu_1, u_1u_2, \dots, u_{k-2}u_{k-1}\}$ . Clearly,  $T' \in T_{n,\Delta}$  and

$$\sum_{uv\in E(T)-S}f(uv)=\sum_{uv\in E(T')-S'}f(uv).$$

By definition

$$\frac{n}{2}GA_2(T) = \sum_{uv \notin S} f(uv) + \sum_{uv \in S} f(uv) = \sum_{uv \in E(T)-S} f(uv) + (k-1)\sqrt{n-1}, \quad (1)$$
and

 $\frac{n}{2}GA_2(T') = \sum_{uv \notin S'} f(uv) + \sum_{uv \in S'} f(uv) = \sum_{uv \in E(T')-S'} f(uv) + \sum_{i=1}^{k-1} \sqrt{i(n-i)}.$  (2) Combining (1), (2) and the fact that  $k \ge 3$ , we obtain  $GA_2(T) < GA_2(T')$ , as desired.

Case 2. *u* is a support vertex.

By Case 1, we may assume that u is not an end-support vertex and  $deg(u_1) = 1$ . Suppose  $deg(u_2) = 2$  and  $T_{u_2}$  is the component of  $T - uu_2$  containing  $u_2$ . Since, by the choice of vertex u, d(u, v) is as large as possible, we may assume that  $T_{u_2}$  is the path  $u_2x_1x_2...x_t, t \ge 1$ . Let T' be the tree obtained from  $T - uu_1$  by adding the pendant edge  $x_tu_1$  to this graph. Let  $S = \{uu_1, uu_2, u_2x_1, x_1x_2, ..., x_{t-1}x_t\}$  and  $S' = \{uu_2, u_2x_1, x_1x_2, ..., x_{t-1}x_t, u_1x_t\}$ . Clearly,  $T' \in T_{n,\Delta}$  and

$$\sum_{uv\in E(T)-S}f(uv)=\sum_{uv\in E(T')-S'}f(uv).$$

By definition

$$\frac{n}{2}GA_2(T) = \sum_{uv \in E(T)-S} f(uv) + \sum_{i=1}^{t+1} \sqrt{i(n-i)} + \sqrt{n-1}, \quad (3)$$

and

$$\frac{n}{2}GA_2(T') = \sum_{uv \in E(T')-S'} f(uv) + \sum_{i=1}^{t+2} \sqrt{i(n-i)}.$$
(4)

By (3), (4) and the fact that  $n \ge t + 4$ , we obtain  $GA_2(T) < GA_2(T')$ .

Case 3. *u* is not a support vertex.

Suppose  $T_{u_1}$  and  $T_{u_2}$  are the components of  $T - \{uu_1, uu_2\}$  containing  $u_1$  and  $u_2$ , respectively. By the choice of vertex u, we may assume that  $T_{u_1} = u_1 x_1 x_2 \dots x_s$ ,  $s \ge 1$  and  $T_{u_2} = u_2 y_1 y_2 \dots y_t$ ,  $t \ge 1$ . Then  $deg(x_i) = deg(y_j) = 2$ ,  $1 \le i \le s - 1$ ,  $1 \le j \le t - 1$ , and  $deg(x_s) = deg(y_t) = 1$ . Let T' be the tree obtained from  $T - T_{u_2}$  by adding the path  $x_s y_t y_{t-1} \dots y_1 u_2$  to this graph. Let

$$S = \{uu_1, u_1x_1, x_1x_2, \dots, x_{s-1}x_s\} \cup \{uu_2, u_2y_1, y_1y_2, \dots, y_{t-1}y_t\},\$$

and

$$S' = \{uu_1, u_1x_1, x_1x_2, \dots, x_{s-1}x_s\} \cup \{x_sy_t, u_2y_1, y_1y_2, \dots, y_{t-1}y_t\}.$$

Clearly, 
$$T' \in T_{n,\Delta}$$
 and

$$\sum_{uv\in E(T)-S}f(uv) = \sum_{uv\in E(T')-S'}f(uv)$$

By definition we have

$$\frac{n}{2}GA_2(T) = \sum_{uv \in E(T)-S} f(uv) + \sum_{i=1}^{S+1} \sqrt{i(n-i)} + \sum_{i=1}^{t+1} \sqrt{i(n-i)}, \quad (5)$$

and

$$\frac{n}{2}GA_2(T') = \sum_{uv \in E(T')-S'} f(uv) + \sum_{i=1}^{s+t+2} \sqrt{i(n-i)}.$$
(6)

Applying Observation 4 and inequalities (5) and (6), we conclude that  $GA_2(T) < GA_2(T')$ . This complete the proof.

A spider is a tree with at most one vertex of degree more than 2, called the center of the spider (if no vertex is of degree more than two, then any vertex can be the center). A leg of a spider is a path from the center to a vertex of degree 1. Thus, a star with k edges is a spider of k legs, each of length 1, and a path is a spider of 1 or 2 legs.

**Lemma 6.** Let T be a spider of order n with  $k \ge 3$  legs. If T has two legs of length at least 2, then there is a spider T' of order n with k legs such that  $GA_2(T) < GA_2(T')$ .

*Proof.* Let v be the center of T and  $N(v) = \{v_1, v_2, ..., v_k\}$ . Root T at v. Assume, without loss of generality, that  $deg(v_1) = deg(v_2) = 2$  and let  $v_1x_1x_2...x_s$  and  $v_2y_1y_2...y_t$  be two legs of T. Let T' be the tree obtained from T be deleting the edges  $x_1x_2,...,x_{s-1}x_s$  and adding the edges  $x_1y_t, x_1x_2,..., x_{s-1}x_s$ . Suppose

$$S = \{vv_1, v_1x_1, x_1x_2, \dots, x_{s-1}x_s\} \cup \{vv_2, v_2y_1, y_1y_2, \dots, y_{t-1}y_t\},\$$

and

$$S' = \{vv_1, y_t x_1, x_1 x_2, \dots, x_{s-1} x_s\} \cup \{vv_2, v_2 y_1, y_1 y_2, \dots, y_{t-1} y_t\}.$$

Clearly

$$\sum_{uv\in E(T)-S}f(uv)=\sum_{uv\in E(T')-S'}f(uv).$$

By definition we have

$$\frac{n}{2}GA_2(T) = \sum_{uv \in E(T)-S} f(uv) + \sum_{i=1}^{S+1} \sqrt{i(n-i)} + \sum_{i=1}^{t+1} \sqrt{i(n-i)}, \quad (7)$$

and

$$\frac{n}{2}GA_2(T') = \sum_{uv \in E(T')-S'} f(uv) + \sum_{i=1}^{s+t+1} \sqrt{i(n-i)} + \sqrt{n-1}.$$
 (8)

By Observation 4, equalities (7) and (8) and the fact that  $n \ge s + t + 4$  we obtain  $GA_2(T) < GA_2(T')$ .

We are now ready to prove the main theorem of this section.

**Theorem 7.** For any tree  $T \in T_{n,\Delta}$  of order  $n \ge 2$ ,

$$GA_2(T) \leq \frac{2}{n} \left( (\Delta - 1)\sqrt{n - 1} + \sum_{i=1}^{n - \Delta} \sqrt{i(n - i)} \right).$$

The equality holds if and only if T is a spider with at most one leg of length at least two.

*Proof.* Let  $T_1$  be a tree of order  $n \ge 2$  with maximum degree  $\Delta$  such that

 $GA_2(T_1) = \max\{GA_2(T) \mid T \text{ is a tree of order } n \text{ with maximum degree } \Delta\}.$ 

Let v be a vertex with maximum degree  $\Delta$ . Root  $T_1$  at v. If  $\Delta = 2$ , then  $T_1$  is a path of order n and the result follows by Theorem A. Let  $\Delta \ge 3$ . By the choice of  $T_1$ , we deduce from Lemma 5 that  $T_1$  is a spider with center v. It follows from Lemma 6 and the choice of  $T_1$ that  $T_1$  has at most one leg of length at least two. First let all legs of  $T_1$  have length one. Then  $T_1$  is a star of order n and the result follows by Theorem B. Now let  $T_1$  have only one leg of length at least two. Then

$$GA_2(T) = \frac{2}{n} \left( (\Delta - 1)\sqrt{n-1} + \sum_{i=1}^{n-\Delta} \sqrt{i(n-i)} \right).$$

This completes the proof.

## **3** UNICYCLIC GRAPHS

A connected graph with precisely one cycle is called a unicyclic graph. Let the set  $\varphi_{n,\Delta,k}$  consist of all unicycle graphs of order n, maximum degree  $\Delta \ge 3$  and grith k, where  $3 \le k \le n$ . Note that if G is a cycle of order n, then  $GA_2(G) = n$ . Let  $G \in \varphi_{n,\Delta,k}$ . In this section we assume that the k-cycle of G is  $C_k = (w_1, w_2, \dots, w_k)$ . In addition for a vertex  $u \in V(C_k)$  we let  $T_u$  be the connected component of  $G \setminus E(C_k)$  containing u. Note that  $T_u$  is a tree and we assume u is the root of this tree. Without loss of generality, we also assume one of the vertices of  $T_{w_1}$ , say v, is of degree  $\Delta$ .

**Lemma 8.** Let  $G \in \varphi_{n,\Delta,k}$  and v be a vertex of maximum degree  $\Delta$ . Let C be the only cycle of G,  $u \in V(C)$  and  $u \neq v$ . If  $T_u$  is a spider with at least two legs, then there is a graph  $G' \in \varphi_{n,\Delta,k}$  such that  $GA_2(G) < GA_2(G')$ .

*Proof.* Assume  $T_u$  has  $\ell$  legs with lengths  $t_1, t_2, ..., t_\ell$  and  $\sum_{i=1}^{\ell} t_i = s$ . Let the graph G' be obtained from  $G \setminus E(T_u)$  by attaching a path  $P_s$  to vertex u. Obviously,  $G' \in \varphi_{n,\Delta,k}$ . A simple calculation shows that

$$GA_{2}(G') - GA_{2}(G) = \frac{2}{n} \left[ \sum_{i=1}^{s} \sqrt{i(n-i)} - \sum_{j=1}^{\ell} \sum_{i=1}^{t_{j}} \sqrt{i(n-i)} \right]$$

Apply Observation 4 to obtain  $GA_2(G') - GA_2(G) > 0$ .

**Lemma 9.** Let  $G \in \varphi_{n,\Delta,k}$  and  $deg(u) \ge 3$ , where  $u \in T_{w_i}$ ,  $u \ne w_i$ , for some  $2 \le i \le k$ . Then there is a graph  $G' \in \varphi_{n,\Delta,k}$  such that  $GA_2(G) < GA_2(G')$ .

*Proof.* Without loss of generality, we may assume u has the largest distance from  $w_i$  among all the vertices of  $T_{w_i}$  whose degree is at least 3. This implies that  $T_u$  is a spider with at least two legs. Let G' be the graph obtained from G by replacing  $T_u$  with a path with the same order as  $T_u$ . A calculation similar to that presented in Lemma 8 shows that  $GA_2(G') - GA_2(G) > 0$ .

**Lemma 10.** Let  $G \in \varphi_{n,\Delta,k}$  and  $T_{w_i}$  and  $T_{w_j}$  be paths of length at least 1 for some  $2 \leq i, j \leq k, i \neq j$ . Then there is a graph  $G' \in \varphi_{n,\Delta,k}$  such that  $GA_2(G) < GA_2(G')$ .

*Proof.* Let  $\ell_1$  and  $\ell_2$  be the length of the paths  $T_{w_i}$  and  $T_{w_j}$ , respectively. Let G' be the graph obtained from G by removing  $T_{w_i}$  and  $T_{w_j}$  and attaching a path of length  $\ell_1 + \ell_2$  to the vertex u. Then as before one can see that  $GA_2(G) < GA_2(G')$ .

**Lemma 11.** Let  $G \in \varphi_{n,\Delta,k}$  and assume the vertices of the cycle  $C_k$  are all of degree two except  $w_1$  and  $w_i$ ,  $i \neq 1$ . If the distance of  $w_i$  from  $w_1$  is not  $\lceil (k-1)/2 \rceil$ , then there is a graph  $G' \in \varphi_{n,\Delta,k}$  such that  $GA_2(G) < GA_2(G')$ .

*Proof.* Let G' be the graph obtained from G by removing  $T_{w_i}$  and attaching it to vertex  $w_j$ , where  $j = \lfloor (k-1)/2 \rfloor$ . Then one can see that  $GA_2(G) < GA_2(G')$ .

Now we consider the graph  $G \in \varphi_{n,\Delta,k}$  with  $deg(w_i) = 2$  for all  $2 \le i \le k, i \ne [(k-1)/2]$  and  $deg(w_j) \ge 2$ , where j = [(k-1)/2]. By Lemma 9, in order to maximize  $GA_2(G)$ ,  $T_v$  must be a spider and  $deg_G(w_1) = 3$  if  $w_1 \ne v$ .

**Lemma 12.** Let  $G \in \varphi_{n,\Delta,k}$  and  $w_1 \neq v$ . Then there is a graph  $G' \in \varphi_{n,\Delta,k}$  such that  $GA_2(G) < GA_2(G')$ .

*Proof.* Let G' be the graph obtained from  $G \setminus T_{w_1}$  by attaching a path of order  $|V(T_{w_1})| - \Delta + 2$  to the end vertex of the path  $T_{w_j}$  which is different from  $w_j$ ,  $j = \lceil (k-1)/2 \rceil$  and adding  $\Delta - 2$  pendant edges at vertex  $w_1$ . Obviously,  $G' \in \varphi_{n,\Delta,k}$  and it is straightforward to verify that  $GA_2(G) < GA_2(G')$ .

By Lammas 8–12 we obtain the following result.

**Corollary 13.** Let  $H \in \varphi_{n,\Delta,k}$  be the graph which consists of a cycle  $C_k = (w_1, w_2, ..., w_k)$ with  $\Delta - 2$  pendant edges at vertex  $w_1$  and a path of order  $n - k - \Delta + 2$  at vertex  $w_j$ , where  $j = \lfloor (k-1)/2 \rfloor$ . Then for every  $G \in \varphi_{n,\Delta,k}$ ,  $GA_2(G) \leq GA_2(H)$ .

We are now ready to state the main theorem of this section.

**Theorem 14.** For any unicycle graph G of order n, girth k and maximum degree  $\Delta \ge 3$ , if k is odd, then

$$\begin{aligned} GA_2(G) &\leq \frac{2}{n} \left( (\Delta - 2)\sqrt{n - 1} + \sum_{i=1}^{n - k - \Delta + 2} \sqrt{i(n - i)} \right) \\ &+ \frac{2(k - 1)}{n - 1} \sqrt{\left(\frac{k - 1}{2} + \Delta - 2\right)(n - \frac{k - 1}{2} - \Delta + 1)} + \frac{2}{\Delta + k - 3} \sqrt{\frac{k - 1}{2}\left(\frac{k - 1}{2} + \Delta - 2\right)} \\ &+ \frac{2}{n - \Delta + 1} \sqrt{\frac{k - 1}{2}(n - \frac{k - 1}{2} - \Delta + 1)}, \end{aligned}$$

and if *k* is even, then

$$GA_{2}(G) \leq \frac{2}{n} \left( (\Delta - 2)\sqrt{n - 1} + \sum_{i=1}^{n - k - \Delta + 2} \sqrt{i(n - i)} + k \sqrt{(\frac{k}{2} + \Delta - 2)(n - \frac{k}{2} - \Delta + 2)} \right).$$

The equality holds if and only if G is the graph H given in Corollary 13.

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